# Multimedia Information Extraction and Retrieval 

## Indexing and Query Answering

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## Recall basic indexing pipeline



## Tokenization

- Input: "Friends, Romans and Countrymen"
- Output: Tokens
- Friends
- Romans
- Countrymen
- Each such token is now a candidate for an index entry, after further processing
- Described below
- But what are valid tokens to emit?


## Tokenization

- Issues in tokenization:
- Finland's capital $\rightarrow$

Finland? Finlands? Finland's?

- Hewlett-Packard $\rightarrow$ Hewlett and Packard as two tokens?
- State-of-the-art: break up hyphenated sequence.
- co-education?
- the hold-him-back-and-drag-him-away-maneuver?
- It's effective to get the user to put in possible hyphens
- San Francisco: one token or two? How do you decide it is one token?


## Numbers

- 3/12/91

Mar. 12, 1991

- 55 B.C.
- B-52
- My PGP key is 324a3df234cb23e
- 100.2.86.144
- Often, don't index as text.
- But often very useful: think about things like looking up error codes/stacktraces on the web
- (One answer is using n-grams: later)
- Will often index "meta-data" separately
- Creation date, format, etc.


## Tokenization: Language issues

- L'ensemble $\rightarrow$ one token or two?
- L? L'? Le ?
- Want l'ensemble to match with un ensemble
- German noun compounds are not segmented
- Lebensversicherungsgesellschaftsangestellter
- 'life insurance company employee’


## Normalization

- Need to "normalize" terms in indexed text as well as query terms into the same form
- We want to match U.S.A. and USA
- We most commonly implicitly define equivalence classes of terms
- e.g., by deleting periods in a term
- Alternative is to do asymmetric expansion:
- Enter: window Search: window, windows
- Enter: windows Search: Windows, windows
- Enter: Windows Search: Windows
- Potentially more powerful, but less efficient


## Normalization: other languages

- Accents: résumé vs. resume.
- Most important criterion:
- How are your users like to write their queries for these words?
- Even in languages that standardly have accents, users often may not type them
- German: Tuebingen vs. Tübingen
- Should be equivalent


## Case folding

- Reduce all letters to lower case
- exception: upper case (in mid-sentence?)
- e.g., General Motors
- Fed vs. fed
- SAIL vs. sail
- Often best to lowercase everything, since users will use lowercase regardless of 'correct' capitalization...


## Stop words

- With a stop list, you exclude from dictionary entirely the commonest words. Intuition:
- They have little semantic content: the, a, and, to, be
- They take a lot of space: ~30\% of postings for top 30
- But the trend is away from doing this:
- Good compression techniques means the space for including stopwords in a system is very small
- Good query optimization techniques mean you pay little at query time for including stop words.
- You need them for:
- Phrase queries: "King of Denmark"
- Various song titles, etc.: "Let it be", "To be or not to be"
- "Relational" queries: "flights to London"


## Thesauri

- Handle synonyms and homonyms
- Hand-constructed equivalence classes
- e.g., car = automobile
- color = colour
- Rewrite to form equivalence classes
- Index such equivalences
- When the document contains automobile, index it under car as well (usually, also vice-versa)
- Or expand query?
- When the query contains automobile, look under car as well


## Lemmatization

- Reduce inflectional/variant forms to base form
- E.g.,
- am, are, is $\rightarrow$ be
- car, cars, car's, cars' $\rightarrow$ car
- the boy's cars are different colors $\rightarrow$ the boy car be different color
- Lemmatization implies doing "proper" reduction to dictionary headword form


## Simpler Form: Stemming

- Reduce terms to their "roots" before indexing
- "Stemming" suggests crude affix chopping
- language dependent
- e.g., automate(s), automatic, automation all reduced to automat.
for example compressed and compression are both accepted as equivalent to compress.
for exampl compress and compress ar both accept as equival to compress


## Porter's Algorithm

- Common algorithm for stemming English
- Results suggest at least as good as other stemming options
- Conventions +5 phases of reductions
- phases applied sequentially
- each phase consists of a set of commands
- sample convention: Of the rules in a compound command, select the one that applies to the longest suffix.


## Porter's Algorithm

- [C](VC)m[V]
- m indicates repetition, $\mathrm{C}=$ consonant, $\mathrm{V}=$ vowel
- X denotes a sequence of Xs
- Examples:
- m=0 TR, EE, TREE, Y, BY
- m=1 TROUBLE, OATS, TREES, IVY
- m=2 TROUBLES, PRIVATE, OATEN
- Conditions:
- *S - the stem ends with S (and similarly for the other letters).
- *v* - the stem contains a vowel.
- *d - the stem ends with a double consonant (e.g. -TT, -SS).
- *o - the stem ends cvc, where the second c is not $\mathrm{W}, \mathrm{X}$ or Y (e.g. -WIL, -HOP).


## Porter's Algorithm

```
Step la
    SSES _> SS
    IES -> I
    SS -> SS
    S ->
```

Step 1b
$(m>0) E E D \rightarrow E E$
(*V*) ED ->
(*V*) ING ->

| caresses | -> caress |  |
| :--- | :--- | :--- |
| ponies | $->$ | $p o n i$ |
| ties | $->$ | ti |
| caress | $->$ | caress |
| cats | $->$ | cat |


| feed | -> feed |  |
| :--- | :--- | :--- |
| agreed | $->$ | agree |
| plastered | -> plaster |  |
| bled | -> bled |  |
| motoring | -> motor |  |
| sing | $->$ | sing |

## Porter's Algorithm

If the second or third of the rules in Step 1 b is successful, the following is done:
$A T->A T E$
$B L->B L E$
$I Z ~->~ I Z E$
(*d and not (*L or *S or *Z))
-> single letter

$$
(m=1 \text { and *O) } \rightarrow E
$$

Step 1c
(*V*) Y -> I
hopp(ing) -> hop
fall(ing) -> fall
hiss(ing) -> hiss
fizz(ed) -> fizz
fail(ing) -> fail
fil(ing) -> file

```
    conflat(ed) -> conflate
    troubl(ed) -> trouble
    siz(ed) -> size
```

    \(\rightarrow\) file
    happy -> happi
sky -> sky

## Porter's Algorithm

| Step 2 |  |  |  |
| :---: | :---: | :---: | :---: |
| $(m>0)$ | ATIONAL | -> | $A T E$ |
| ( $m>0$ ) | TIONAL | -> | TION |
| $(m>0)$ | ENCI | -> | ENCE |
| $(m>0)$ | $A N C I$ | -> | ANCE |
| $(m>0)$ | $I Z E R$ | -> | IZE |
| $(m>0)$ | $A B L I$ | -> | $A B L E$ |
| $(m>0)$ | $A L L I$ | -> | $A L$ |
| $(m>0)$ | ENTLI | -> | ENT |
| $(m>0)$ | $E L I$ | -> | $E$ |
| $(m>0)$ | OUSLI | -> | OUS |
| $(m>0)$ | IZATION | -> | IZE |
| ( $m>0$ ) | $A T I O N$ | -> | $A T E$ |
| $(m>0)$ | $A T O R$ | -> | ATE |
| $(m>0)$ | $A L I S M$ | -> | $A L$ |
| $(m>0)$ | IVENESS | -> | IVE |
| $(m>0)$ | FULNESS | -> | FUL |
| $(m>0)$ | OUSNESS | -> | OUS |
| $(m>0)$ | $A L I T I$ | -> | $A L$ |
| $(m>0)$ | IVITI | -> | IVE |
| ( $m>0$ ) | BILITI | -> | $B L E$ |


| relational | -> | relate |
| :---: | :---: | :---: |
| conditional | -> | condition |
| rational | -> | rational |
| lenci | -> | valence |
| hesitanci | -> | hesitance |
| digitizer |  | ize |
| conformabli | -> | onformable |
| radicalli | -> | adical |
| differentli | -> | different |
| vileli | - $>$ | ile |
| analogousli | -> | analogous |
| vietnamization |  | vietnamize |
| predication | -> | predicate |
| operator | -> | operate |
| feudalism | -> | feudal |
| decisiveness | -> | decisive |
| hopefulness | -> | hopeful |
| callousness | -> | callous |
| formaliti | -> | formal |
| sensitiviti | -> | ensitive |
| sensibiliti | -> | sensible |

## Porter's Algorithm

## Step 3

| $(m>0)$ | ICATE | $->$ | $I C$ |
| :--- | :--- | :--- | :--- |
| $(m>0)$ | $A T I V E$ | $->$ |  |
| $(m>0)$ | $A L I Z E$ | $->$ | $A L$ |
| $(m>0)$ | $I C I T I$ | $->$ | $I C$ |
| $(m>0)$ | $I C A L$ | $->$ | $I C$ |
| $(m>0)$ | $F U L$ | $->$ |  |

Step 4
(m>1) AL ->
(m>1) ANCE ->
(m>1) ENCE ->
(m>1) ER ->
(m>1) IC ->
( $m>1$ ) $A B L E$->
( $m>1$ ) IBLE ->
( $\mathrm{m}>1$ ) ANT ->
(m>1) EMENT ->
(m>1) MENT ->
(m>1) ENT ->
( $m>1$ and (*S or *T)) ION ->
( $m>1$ ) OU ->
(m>1) ISM ->
( $m>1$ ) ATE ->
(m>1) ITI ->
(m>1) OUS ->
(m>1) IVE ->
(m>1) IZE ->

| triplicate | -> | triplic |
| :---: | :---: | :---: |
| formative | -> | form |
| formalize | -> | formal |
| electriciti | -> | electric |
| electrical | -> | electric |
| hopeful | -> | hope |
| goodness | -> | good |
| revival | -> | reviv |
| allowance | -> | allow |
| inference | -> | infer |
| airliner | -> | airlin |
| gyroscopic | -> | gyroscop |
| adjustable | -> | adjust |
| defensible | -> | defens |
| irritant | -> | irrit |
| replacement | -> | replac |
| adjustment | -> | adjust |
| dependent | -> | depend |
| adoption | -> | adopt |
| homologou | -> | homolog |
| communism | -> | commun |
| activate | -> | activ |
| angulariti | -> | angular |
| homologous | -> | homolog |
| effective | -> | effect |
| bowdlerize | -> | bowdler |

## Porter's Algorithm

```
Step 5a
    (m>1) E ->
    (m=1 and not *O) E ->
    rate
    cease
    -> probat
Step 5b
(m > 1 and *d and *L) -> single letter
controll
roll
-> control
    -> roll
```


## Faster postings merges: Skip pointers

## Recall basic merge

- Walk through the two postings simultaneously, in time linear in the total number of postings entries


If the list lengths are $m$ and $n$, the merge takes $O(m+n)$ operations.

Can we do better?
Yes, if index isn't changing too fast.

## Augment postings with (at indexing time)



- Why?
- To skip postings that will not figure in the search results.
- How?
- Where do we place skip pointers?


## Query processing with



Suppose we've stepped through the lists until we process 8 on each list.

When we get to 16 on the top list, we see that its successor is 32.
But the skip successor of 8 on the lower list is 31, so we can skip ahead past the intervening postings.

## Where do we place skips?

- Tradeoff:
- More skips $\rightarrow$ shorter skip spans $\Rightarrow$ more likely to skip. But lots of comparisons to skip pointers.
- Fewer skips $\rightarrow$ few pointer comparison, but then long skip spans $\Rightarrow$ few successful skips.



## Placing skips

- Simple heuristic: for postings of length $L$, use $\sqrt{L}$ evenly-spaced skip pointers.
- This ignores the distribution of query terms.
- Easy if the index is relatively static; harder if $L$ keeps changing because of updates.
- This definitely used to help; with modern hardware it may not
- The cost of loading a bigger postings list outweighs the gain from quicker in memory merging


## Phrase queries

## Phrase queries

- Want to answer queries such as "stanford university" - as a phrase
- Thus the sentence "I went to university at Stanford" is not a match.
- The concept of phrase queries has proven easily understood by users; about $10 \%$ of web queries are phrase queries
- No longer suffices to store only
<term : docs> entries


## A first attempt: Biword indexes

- Index every consecutive pair of terms in the text as a phrase
- For example the text "Friends, Romans, Countrymen" would generate the biwords
- friends romans
- romans countrymen
- Each of these biwords is now a dictionary term
- Two-word phrase query-processing is now immediate.


## Longer phrase queries

- Longer phrases are processed as follows:
- stanford university palo alto can be broken into the Boolean query on biwords:
stanford university AND university palo AND palo alto

Without the docs, we cannot verify that the docs matching the above Boolean query do contain the phrase.

Can have false positives!

## Extended biwords

- Parse the indexed text and perform part-of-speechtagging (POST).
- Bucket the terms into (say) Nouns (N) and articles/ prepositions (X).
- Now deem any string of terms of the form NX*N to be an extended biword.
- Each such extended biword is now made a term in the dictionary.
- Example: catcher in the rye $\mathrm{N} \quad \mathrm{X} \quad \mathrm{X}$
- Query processing: parse it into N's and X's
- Segment query into enhanced biwords
- Look up index


## Issues for biword indexes

- False positives, as noted before
- Index blowup due to bigger dictionary
- For extended biword index, parsing longer queries into conjunctions:
- E.g., the query tangerine trees and marmalade skies is parsed into
- tangerine trees AND trees and marmalade AND marmalade skies
- No standard solution (for all biwords)


## Solution 2: Positional indexes

- Store, for each term, entries of the form:
<number of docs containing term;
doc1: position1, position2 ... ;
doc2: position1, position2 ... ; etc.>


## Positional index example

<be: 993427;
1: 7, 18, 33, 72, 86, 231;
2: 3, 149;
4: 17, 191, 291, 430, 434;
5: $363,367, \ldots>$


- Can compress position values/ offsets
- Nevertheless, this expands postings storage substantially


## Processing a phrase query

- Extract inverted index entries for each distinct term: to, be, or, not.
- Merge their doc:position lists to enumerate all positions with "to be or not to be".
- to:

$$
\begin{aligned}
& \text { 2:1,17,74,222,551; 4:8,16,190,429,433; } \\
& \quad 7: 13,23,191 ; \ldots
\end{aligned}
$$

-be:

- $1: 17,19 ; 4: 17,191,291,430,434 ; 5: 14,19,101 ; \ldots$
- Same general method for proximity searches


## Proximity queries

- LIMIT! /3 STATUTE /3 FEDERAL /2 TORT Here, /k means "within $k$ words of".
- Clearly, positional indexes can be used for such queries; biword indexes cannot.
- Exercise: Adapt the linear merge of postings to handle proximity queries. Can you make it work for any value of $k$ ?


## Positional index size

- You can compress position values/offsets:
- Nevertheless, a positional index expands postings storage substantially
- Nevertheless, it is now standardly used because of the power and usefulness of phrase and proximity queries ... whether used explicitly or implicitly in a ranking retrieval system.


## Positional index size

- Need an entry for each occurrence, not just once per document
- Index size depends on average document why? size
- Average web page has <1000 terms
- SEC filings, books, even some epic poems ... easily 100,000 terms
- Consider a term with frequency $0.1 \%$

| Document size | Postings | Positional postings |
| ---: | ---: | ---: |
| 1000 | 1 | 1 |
| 100,000 | 1 | 100 |

## Rules of thumb

- A positional index is 2-4 as large as a non-positional index
- Positional index size 35-50\% of volume of original text
- Caveat: all of this holds for "Englishlike" languages


## Wild-card queries: *

- mon*: find all docs containing any word beginning "mon".
- Easy with binary tree (or B-tree) lexicon: retrieve all words in range: $\boldsymbol{m o n} \leq \boldsymbol{w}<$ moo
- "mon: find words ending in "mon": harder
- Maintain an additional B-tree for terms backwards.
Can retrieve all words in range: nom $\leq \boldsymbol{w}<\boldsymbol{n o n}$.
Exercise: from this, how can we enumerate all terms meeting the wild-card query pro*cent?


## B-tree

- Binary tree data structure
- Optimized for page-oriented storage of data on harddisks
- Original version: B-tree: R. Bayer and E.M. McCreight. Organization and Maintenance of Large Ordered Indexes. Acta Informatica, vol. 1, no. 3, September 1972.
- leaf nodes are, generally, not in sequential order on disk,
- leaves are connected to form a double-linked list: ${ }^{2}$



## B-tree: Central idea by example

## Insert: Examples (Insert without Split)



Insert new entry with key 4222.
$\rightarrow$ Enough space in node 3, simply insert.
$\rightarrow$ Keep entries sorted within nodes.

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## Query processing

- At this point, we have an enumeration of all those terms in the dictionary that match the wild-card query.
- We still have to look up the postings for each enumerated term.
- E.g., consider the query: se*ate AND fil*er
This may result in the execution of many Boolean $A N D$ queries.


## B-trees handle *'s at the end of a query term

- How can we handle *'s in the middle of query term?
*(Especially multiple *'s)
- The solution: transform every wildcard query so that the *'s occur at the end
- This gives rise to the Permuterm Index.


## Permuterm index

- For term hello index under:
- hello\$, ello\$h, llo\$he, lo\$hel, o\$hell, \$hello where $\$$ is a special symbol.
- Queries:
- X lookup on X\$
* *X lookup on $\mathbf{X} \$$ *
- X*Y lookup on Y\$X*

Exercise!

X* lookup on $\$ \mathbf{X}^{*}$

* $\mathbf{X}^{*}$ lookup on $\mathrm{X}^{*}$ $\mathbf{X} * \mathbf{Y} * \mathbf{Z} \quad$ ???



## Permuterm query processing

- Rotate query wild-card to the right
- Now use B-tree lookup as before.
- Permuterm problem: $\approx$ quadruples lexicon size

Empirical observation for English.

## Bigram indexes

- Enumerate all $k$-grams (sequence of $k$ chars) occurring in any term
- e.g., from text "April is the cruelest month" we get the 2-grams (bigrams)

$$
\begin{aligned}
& \text { \$a, ap,pr,ri, il,l\$,\$i,is,s\$,\$t,th,he,e\$,\$c,cr,ru, } \\
& \text { ue,el,le,es,st,t\$, \$m,mo,on,nt,h\$ }
\end{aligned}
$$

- \$ is a special word boundary symbol
- Maintain an "inverted" index from bigrams to dictionary terms that match each bigram.


## Bigram index example



## Processing $n$-gram wild-cards

- Query mon* can now be run as
- \$m AND mo AND on
- Fast, space efficient.
- Gets terms that match the AND-version of our wildcard query.
- But we'd enumerate moon.
- Must post-filter these terms against query.
- Surviving enumerated terms are then looked up in the term-document inverted index.

