## Scalable RDF Data Management

## ... with a Touch of Uncertainty

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Joint work with:

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## Jeffrey Ullman

Information Extraction
From Wikipedia, the free encyclopedia

Jeffrey David Ullman (born November 22, 1942) is a renowned computer scientist. His textbooks on compilers (various editions are popularly known as the Dragon Book), theory of computation (also known as the Cinderella book), data structures, and databases are regarded as standards in their fields.

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Early life \& Career [edit]

Ullman received a Bachelor of Science degree in Engineering Mathematics from Columbia University in 1963 and his Ph.D. in Electrical Engineering from Princeton University in 1966. He then worked for several years at Bell Labs. From 1969 to 1979 he was a professor at Princeton. Since 1979 he has been a professor at Stanford

## Jeffrey D. Ullman

Born November 22, 1942 (age 69)
Citizenship American
Nationality American
Institutions Stanford University
Alma $\quad$ Columbia University,
mater
Doctoral
advisor
Doctoral
students
Princeton University
Arthur Bernstein, Archie McKellar
Surajit Chaudhuri, Kevin Karplus, David Maier, Harry Mairson, Alberto O. Mendelzon, Jeffrey F. Naughton, Anand Rajaraman, Yehoshua Sagiv,
Mihalis Yannakakis
Known for
database theory, database systems, formal language theory

## Notable

 awardsFellow of the Association for Computing Machinery (1994), ACM SIGMOD Contributions Award (1996), ACM SIGMOD Best Paper Award (1996), Karl V. Karlstrom outstanding educator award (1998), Knuth Prize (2000),
ACM SIGMOD Edgar F. Codd Innovations Award (2006), ACM SIGMOD Test of Time Award (2006), IEEE John von Neumann Medal (2010) University, where he is currently the Stanford W. Ascherman Professor of Computer Science (Emeritus). In 1995 he was inducted as a Fellow of the Association for Computing Machinery and in 2000 he was awarded the Knuth Prize. Ullman is also the co-recipient (with John Hopcroft) of the 2010 IEEE John von Neumann Medal, "For laying the foundations for the fields of automata and language theory and many seminal contributions to theoretical computer science." $[1]$

Ullman's research interests include database theory, data integration, data mining, and education using the information infrastructure. He is one of the founders of the field of database theory, and was the doctoral advisor of an entire generation of students who later became leading database theorists in their own right. He was the Ph.D. advisor of Sergey Brin, one of the co-founders of Google, and served on Google's technical advisory board. He is currently the CEO of Gradiance.

## Books

[edit]

## DBpedia/YAGO et al.

> bornOn(Jeff, 09/22/42) gradFrom(Jeff, Columbia) gradFrom(Jeff, Princeton) hasAdvisor(Jeff, Arthur) hasAdvisor(Surajit, Jeff) knownFor(Jeff, Theory)

>120 M facts for YAGO3 (from Wikipedia infoboxes)

## New fact candidates

author(Jeff, Drag_Book) [0.6] author(Jeff, Cind_Book) [0.8] worksAt(Jeff, Bell_Labs) [0.5] hasAdvisor(Sergej, Jeff) [0.7] type(Jeff, ACM_Fellow) [0.5] type(Jeff, CEO) [0.3]
>100's M additional facts
(from Wikipedia free-text)

- Database Systems: The Complete Book (with H. Garcia-Molina and J. Widom), Prentice-Hall, Englewood Cliffs, NJ,


## Linked-Open-Data Cloud

| Legend |
| :--- |
| Cross Domain |
| Geography |
| Government |
| Life Sciences |
| Linguistics |
| Media |
| Publications |
| Social Networking |
| User Generated |
| Incoming Links |



## Wolfram Alpha

## The "Computational Knowledge Engine"

- Fully implemented in Wolfram-Mathematica
- 10 trillion+ facts
- 50,000+ algorithms and statistical analyses
- 5,000+ templates for visualization and layouts
- 1,000+ domain-specific linguistic analyses

January 24: California Gold Rush
February 2: Treaty of Guadalupe Hidalgo
February 21: Communist Manifesto published
July 4: Marx and Engels publish their "Communist Manifesto"
December 2: Franz Josef I becomes Emperor of Austria and King of Hungary

## Calendar

| January |  |  |  |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Su Mo Tu We Th Fr Sa <br> 2 3 4 5 6 7 8 <br> 9 10 11 12 13 14 15 <br> 16 17 18 19 20 21 22 <br> 23 24 25 26 27 28 29 <br> 30 31      |  |  |  |  |  |  |



## IBM Watson: Deep Question Answering

- William Wilkinson's "An Account of the Principalities of Wallachia and Moldavia" inspired this author's most famous novel
- This town is known as "Sin City" \& its downtown is "Glitter Gulch"
- As of 2010, this is the only former Yugoslav republic in the EU
- 99 cents got me a 4-pack of Ytterlig coasters from this Swedish chain
- U.S. City: largest airport is named for a World War II Hero; its second largest for a World War II Battle


## Question

 classification \& decomposition
D. Ferrucci et al.: Building Watson: An Overview of the DeepQA Project. Al Magazine, Fall 2010.

WIKIPEDIA The Free Encyclopedia

freebase

## RDF-Centered Research Topics

## Information Extraction

[SIGMOD'09, WebDB'10, PODS'10, WSDM'11, CIKM'12, CLEF/INEX'11/'12, LDOW'14, TACL'16]

- Uncertain RDF Data \& Probabilistic Databases
[ICDE'08, VLDB-J'08, SSDBM'10, BTW'11, CIKM'11, ICDE'13, PVLDB'14, VLDB PhD Workshop'15]
- Scalable RDF Indexing \& SPARQL Query Processing
[SIGMOD'14, SWIM'14, SIGMOD'16]


## Named-Entity Recognition \& Disambiguation

$$
\begin{aligned}
& \text { "David played for manu, real, and la galaxy. } \\
& \text { His wife posh performed with the spice girls." }
\end{aligned}
$$

- State-of-art approaches recognize named entities and then disambiguate these entities in two strictly separated phases.


## Named-Entity Recognition \& Disambiguation



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## Named-Entity Recognition \& Disambiguation



- State-of-art approaches recognize named entities and then disambiguate these entities in two strictly separated phases.


## Named-Entity Recognition \& Disambiguation

"David played for manu, real, and la galaxy.

His wife posh performed with the spice girls."


- State-of-art approaches recognize named entities and then disambiguate these entities in two strictly separated phases.


## Joint Named-Entity Recognition \& Disambiguation

"David played for manu, real, and la galaxy.
His wife posh performed with the spice girls."


- J-NERD jointly recognizes and disambiguates named entities with respect to a background knowledge base such as YAGO.


## Conditional Random Field in J-NERD



## Conditional Random Field in J-NERD



## Conditional Random Field in J-NERD



## Conditional Random Field in J-NERD



- Probability distribution over possible tokens $x$ and combined NER/D labels $y$

$$
p(\mathbf{x}, \mathbf{y})=\frac{1}{Z} \prod_{A} \mathcal{F}_{A}\left(\mathbf{x}_{A}, \mathbf{y}_{A}\right)
$$

- Probabilistic inference: find the most likely $\mathbf{y}^{*}=\arg \max _{\mathbf{y}} p(\mathbf{y} \mid \mathbf{x})$ labels $\boldsymbol{y}$, given the observed tokens $\boldsymbol{x}$
- Viterbi algorithm (dynamic programming) for fast and exact inference


## Conditional Random Field in J-NERD

## CRF with cross-sentence

# [Nguyen,Theobald,Weikum: LDOW'14, 

 dependencies:

- Probability distribution over possible tokens $x$ and combined NER/D labels $y$

$$
p(\mathbf{x}, \mathbf{y})=\frac{1}{Z} \prod_{A} \mathcal{F}_{A}\left(\mathbf{x}_{A}, \mathbf{y}_{A}\right)
$$

- Probabilistic inference: find the most likely $\mathbf{y}^{*}=\arg \max _{\mathbf{y}} p(\mathbf{y} \mid \mathbf{x})$ labels $\boldsymbol{y}$, given the observed tokens $\boldsymbol{x}$
- General factor graphs: MCMC-style sampling for approximate inference


## Conditional Random Field in J-NERD

## CRF with cross-sentence

[Nguyen,Theobald,Weikum: LDOW'14, dependencies:

Evaluation on the
CoNLL newswire
collection with
YAGO2 ground-
truth annotations
$(1,244$ labeled articles $)$

| Method | Prec | Rec | F $_{1}$ |
| :---: | :---: | :---: | :---: |
| $P$-NERD | 80.1 | 75.1 | 77.5 |
| $J$-NERD | $\mathbf{8 1 . 9}$ | 75.8 | $\mathbf{7 8 . 7}$ |
| AIDA-light | 78.7 | $\mathbf{7 6 . 1}$ | 77.3 |
| TagMe | 64.6 | 43.2 | 51.8 |
| SpotLight | 71.1 | 47.9 | 57.3 |

## Ultimate PhD Challenge (I)



- All of the current NED tools (incl. AIDA, J-NERD, Spotlight, TagMe) get this sentence wrong!
- Humans (usually) get it right, though.


## RDF-Centered Research Topics

- Information Extraction
[SIGMOD'09, WebDB'10, PODS'10, WSDM'11, CIKM'12, CLEF/INEX'11/'12, LDOW'14, TACL'16]

Uncertain RDF Data \& Probabilistic Databases
[ICDE'08, VLDB-J'08, SSDBM'10, BTW'11, CIKM'11, ICDE'13, PVLDB'14, VLDB PhD Workshop'15]

- Scalable RDF Indexing \& SPARQL Query Processing
[SIGMOD'14, SWIM'14, SIGMOD'16]


## Probabilistic Database

A probabilistic database $\mathbf{D}^{p}$ (compactly) encodes a probability distribution over a finite set of deterministic database instances $\mathbf{D}_{i}$.


- Special Cases:
(II) $\mathrm{D}^{p}$ tuple-independent

| worksAt(sub, obj) |  | $p$ |
| :---: | :---: | :--- |
| Jeff | Stanford | 0.6 |
| Jeff | Princeton | 0.7 |

(III) D ${ }^{p}$ block-independent

| worksAt(sub, obj) |  | p |
| :---: | :---: | :--- |
| Jeff | Stanford | 0.6 |
|  | Princeton | 0.4 |

Note:
(I) and (II) here are equivalent; (II) and (III) not!

Query Answering Problem: ("Marginal Probabilities" of Query Answers) Run query $\mathbf{Q}$ against each instance $\mathbf{D}_{i}$; for each answer tuple $t_{j}$, $\mathrm{P}\left(t_{j}\right)$ is the sum of the probabilities of all instances $\mathbf{D}_{i}$ where $t_{j}$ exists.

## Flashback: Stanford Trio System

## Uncertainty-Lineage Databases (ULDBs)

1. Alternatives
2. '?' (Maybe) Annotations
3. Confidence values
4. Lineage

## Trio's Data Model

## 1. Alternatives: uncertainty about value



## Trio's Data Model

## 1. Alternatives <br> 2. '?' (Maybe): uncertainty about presence



## Trio's Data Model

## 1. Alternatives

2. '?' (Maybe) Annotations
3. Confidences: weighted uncertainty

| Saw (witness, color, car) |  |
| :---: | :---: |
| Amy | red, Honda 0.5 \|| red, Toyota 0.3 || orange, Mazda 0.2 |
| Betty | blue, Acura 0.6 |

Still six possible instances, but each with a probability

## So Far: Data Model is Not Closed

| Saw (witness, car) |  |
| :---: | :---: |
| Cathy | Honda \|| Mazda |
|  | Jimmy, Toyota \|| Jimmy, Mazda |
| Billy, Honda \|| Frank, Honda |  |
| Hank, Honda |  |

## Suspects $=\Pi_{\text {person }}($ Saw $\bowtie$ Drives $)$

| Suspects |  |  |
| :---: | :---: | :---: |
| Jimmy | $\boldsymbol{?}$ | CANNOT correctly |
| Billy $\boldsymbol{\\|}$ Frank | $\boldsymbol{?}$ | capture possible |
| Hank | $\boldsymbol{?}$ | result instances |

## Example with Lineage

| ID | Saw (witness, car) |  |
| :---: | :---: | :---: |
| 11 | Cathy | Honda \|| Mazda |


| ID | Drives (person, car) |
| :---: | :---: |
| 21 | Jimmy, Toyota II Jimmy, Mazda |
| 22 | Billy, Honda \|| Frank, Honda |
| 23 | Hank, Honda |

Suspects $=T$ person $($ Saww $\bowtie$ Drives $)$

$$
\begin{aligned}
& \lambda(31)=(11,2) \wedge(21,2) \\
& \lambda(32,1)=(11,1) \wedge(22,1) ; \lambda(32,2)=(11,1) \wedge(22,2) \\
& \lambda(33)=(11,1) \wedge 23
\end{aligned}
$$

## Example with Lineage

| ID | Saw (witness, car) |  |
| :---: | :---: | :---: |
| 11 | Cathy | Honda II Mazda |


| ID | Drives (person, car) |
| :---: | :---: |
| 21 | Jimmy, Toyota \|| Jimmy, Mazda |
| 22 | Billy, Honda \|| Frank, Honda |
| 23 | Hank, Honda |

Suspects $=\Pi_{\text {person }}($ Saw $\bowtie$ Drives $)$

| ID | Suspects |  |  |
| :---: | :---: | :---: | :---: |
| 31 | Jimmy | $?$ | $\lambda(31)=(11,2) \wedge(21,2)$ |
| 32 | Billy $\\|$ Frank | $?$ | $\lambda(32,1)=(11,1) \wedge(22,1) ; \lambda(32,2)=(11,1) \wedge(22,2)$ |
| 33 | Hank | $?$ | $\lambda(33)=(11,1) \wedge 23$ |



Correctly captures the possible result instances

## Operational Semantics



Completeness: any (finite) set of possible instances can be represented
(will be coming back to this subtlety again later...)

## Summary on Trio's Data Model

## Uncertainty-Lineage Databases (ULDBs)

1. Alternatives
2. '?' (Maybe) Annotations
3. Confidence values
4. Lineage

Theorem: ULDBs are closed and complete.
Formally studied properties like minimization, equivalence, approximation and membership based on lineage. [Benjelloun, Das Sarma, Halevy, Widom, Theobald: VLDB-J. 2008]

W Barack Obama - Wikipedia, the free enc..
< - https://en.wikipedia.org/wiki/Obama
Donate to Wikipedia Barack Hussein

- Interaction Help About Wikipedia Community portal Recent changes Contact page
> Toolbox
- Print/export
- Languages

Acèh
Afrikaans
Alemannisch
Kの7Cร
Fnglisc
Апऽсшәа
(1) $x_{2}$

Aragonés
risire
Asturianu
Avañe'é
Aвap
Aymar aru
Azərbaycanca
Bamanankan
বাংলা
Bahasa Banjar
Bân-lâm-gú
Basa Banyumasan
Башкортса
Беларуская
Беларуская
(тарашкевіца)
भोजपुरी
Bikol Central
Bislama
Български

## Obama II

Barack Obama


44th President of the United States Incumbent
Assumed office
January 20, 2009
Vice President Joe Biden
Preceded by George W. Bush
United States Senator

## from Illinois

## In office

January 3, 2005 - November 16, 2008
Preceded by Peter Fitzgerald
Succeeded by Roland Burris
Member of the Illinois Senate
from the 13th District

## In office

January 8, 1997 - November 4, 2004

## Preceded by Alice Palmer

Succeeded by Kwame Raoul

## Personal details

 August 4, 1961 (age 52) Honolulu, Hawaii, U.S.
## back to Wikipedia

| fox |  |  |  |  |  |  |  | X |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| rack Obama citizenship conspiracy th... $\quad+$ |  |  |  |  |  |  |  |  |
| Q https://en.wikipedia.org/wiki/Barack_c $\leftrightarrows$ C |  |  |  | Google | , |  |  | * |

## Born in Kenya [ edit source | edit beta]

Some opponents of Obama's presidential eligibility claim that he was born in Kenve, and was therefore not born a United States citizen.
Whether Obama having been born outside the U.S. would have invalidated his U.S. citizenship at birth is debated. Andrew Malcolm, of the Los Angeles Times, has argued that Obama would still be eligible for the presidency, irrespective of where he was born, because his mother was an American citizen, saying that Obama's mother "could have been on Mars when wee Barry emerged and he'd still be American. ${ }^{[59]}$ A contrary view is promoted by UCLA Law Professor Eugene Volokh, who has said that in the hypothetical scenario that Obama was born outside the U.S., he would not be a natural-born citizen, since the then-applicable law would have required Obama's mother to have been in the U.S. at least "five years after the age of 14 ". but Ann Dunham was three months shy of her 19 th birthday when Obama was born. ${ }^{[60]}$

## Obama's paternal step-grandmother's version of events [ edit source | edit beta ]

An incorrect but popularly reported claim is that his father's stepmother, Sarah Obama, told Anabaptist Bishop Ron McRae in a recorded transatlantic telephone conversation that she was present when Obama was born in Kenya.

## bornln(Barack, Hawaii)

bemininidanack kening)

## Soft Rules vs. Hard Rules

(Soft) Deduction Rules vs.
(Hard) Consistency Constraints

- People may live in more than one

Deductive Database: Datalog, Core of SQL \& Relational Algebra, RDF/S, OWL2-RL, etc. livesIn $(x, y) \Leftarrow$ marriedTo( $x, z) \wedge$ liv. livesIn $(x, y) \Leftarrow$ hasChild $(x, z) \wedge$ livesIn( $z, y$ ) [0.5]

- People are not born in different placonlan diffnment datan bornIn $(x, y) \wedge \operatorname{bornIn}(x, z) \Rightarrow y=z \quad$ More General FOL bornon $(x, y) \wedge \operatorname{bornOn}(x, z) \Rightarrow y=z \quad$ Constraints:

Datalog plus constraints, owl:FunctionalProperty, owl:disjointWith, etc.

- People are not married to more tha (at the same time, in most countries?) marriedTo( $x, y, t_{1}$ ) ^ marriedTo( $x, z, t_{2}, \wedge$ ymi

$$
\Rightarrow \operatorname{disjoint}\left(\mathrm{t}_{1}, \mathrm{t}_{2}\right)
$$

Deductive Grounding w/ Lineage
(SLD Resolution in Datalog/Prolog)
[Yahya,Theobald: RuleML'11, Dylla,Miliaraki,Theobald: ICDE'13]


## Rules

```
hasAdvisor(x,y) ^
worksAt(y,z)
    => graduatedFrom(x,z)
graduatedFrom(x,y) ^
graduatedFrom(x,z)
    y=z
```


## Base Facts

| graduatedFrom(Surajit, Princeton) [0.7] |
| :--- |
| graduatedFrom(Surajit, Stanford) [0.6] |
| graduatedFrom(David, Princeton) [0.9] |
| hasAdvisor(Surajit, Jeff) [0.8] |
| hasAdvisor(David, Jeff) [0.7] |
| worksAt(Jeff, Stanford) [0.9] |
| type(Princeton, University) [1.0] |
| type(Stanford, University) [1.0] |
| type(Jeff, Computer_Scientist) [1.0] |
| type(Surajit, Computer_Scientist) [1.0] |
| type(David, Computer_Scientist) [1.0] |

graduatedFrom(Surajit, Stanford) [0.6] graduatedFrom(David, Princeton) [0.9] hasAdvisor(Surajit, Jeff) [0.8] hasAdvisor(David, Jeff) [0.7]
worksAt(Jeff, Stanford) [0.9] type(Princeton, University) [1.0] type(Stanford, University) [1.0] type(Jeff, Computer_Scientist) [1.0] type(Surajit, Computer_Scientist) [1.0] type(David, Computer_Scientist) [1.0]

## Lineage \& Possible Worlds

## Query graduatedFrom(Surajit, y)

| $0.7 \times(1-0.888)=0.078$ |
| :---: |
| graduatedFrom |
| (Surajit, |
| $\mathbf{Q}_{1}$ Princeton) |

(1-0.7) $\times 0.888=0.266$
raduatedFrom
(Surajit,
$A \wedge \neg(B \vee(C \wedge D))$

[Das Sarma,Theobald,Widom: ICDE’08,
Dylla,Miliaraki,Theobald: ICDE'13]

1) Deductive Grounding

- Top-down Datalog evaluation
- Plus tracing the lineage of individual query answers

2) Lineage DAGs

- Grounded soft \& hard rules
- Base facts with confidences

3) Probabilistic Inference
$\rightarrow$ Compute marginals:
$P(Q)$ : sum up the probabilities of all possible worlds that entail the query answers
$P(Q \mid H)$ : drop "impossible worlds"

## 

$P\left(Q_{1}\right)=0.0784 \quad P\left(Q_{1} \mid H\right)=0.0784 / 0.412$
$P\left(Q_{2}\right)=0.2664$ $=0.1903$
$P\left(Q_{2} \mid H\right)=0.2664 / 0.412$ $=0.6466$

| A:0.7 | B:0.6 | C:0.8 | D:0.9 | $\begin{aligned} & Q_{2}: \\ & \neg A \wedge(B \vee(C \wedge D)) \end{aligned}$ | $\mathbf{P}(\mathbf{W})$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 1 | 0 | $0.7 \times 0.6 \times 0.8 \times 0.9=0.3024$ |
| 4 | 1 | 1 | 0 | 0 | $0.7 \times 0.6 \times 0.8 \times 0.1=0.0336$ |
| 1 | 1 | 0 | 1 | 0 | $\ldots=0.0756$ |
| 1 | 1 | 0 | 0 | 0 | $\ldots=0.0084$ |
| 1 | 0 | 1 | 1 | 0 | $\ldots=0.2016$ |
| 1 | 0 | 1 | 0 | 0 | $\ldots=0.0224$ |
| 1 | 0 | 0 | 1 | 0 | $\ldots=0.0504$ |
| 1 | 0 | 0 | 0 | 0 | $\ldots=0.0056$ |
| 0 | 1 | 1 | 1 | 1 | $0.3 \times 0.6 \times 0.8 \times 0.9=0.1296$ |
| 0 | 1 | 1 | 0 | 1 | $0.3 \times 0.6 \times 0.8 \times 0.1=0.0144$ |
| 0 | 1 | 0 | 1 | 1 | $0.3 \times 0.6 \times 0.2 \times 0.9=0.0324$ |
| 0 | 1 | 0 | 0 | 1 | $0.3 \times 0.6 \times 0.2 \times 0.1=0.0036$ |
| 0 | 0 | 1 | 1 | 1 | $0.3 \times 0.4 \times 0.8 \times 0.9=0.0864$ |
| 0 | 0 | 1 | 0 | 0 | $\ldots=0.0096$ |
| 0 | 0 | 0 | 1 | 0 | $\ldots=0.0216$ |
| 0 | 0 | Hard rule $H: \neg A \vee \neg(B \vee(C \wedge D))$ |  |  |  |

## Dichotomy of Queries

[Suciu \& Dalvi: SIGMOD’05 Tutorial on "Foundations of Probabilistic Answers to Queries"]
A probabilistic database $\mathrm{D}^{p}$ (compactly) encodes a probability distribution over a finite set of deterministic database instances $\mathbf{D}_{i}$.

Is there any professor who works at a university that is located in CA?
Q() :- isProfessor(pers), worksAt(pers,uni), located(uni, CA)


Theorem: The query answering problem for the above join query over a tuple-independent probabilistic database is \#P-hard.

## Inference in Probabilistic Databases

## Safe query plans [Dalvi \& Suciu: VLDB-J'07+J-ACM'12]

- Can propagate confidences along with relational operators.

Read-once functions [Sen et al.: PVLDB'10; Olteanu \& Huang: SUM'08]

- Can factorize Boolean formula (in polynomial time) into read-once form, where every variable occurs at most once.

Knowledge compilation [Olteanu et al.: ICDE'10; ICDT'11; VLDB- $\mathrm{J}^{\prime} 13$ ]

- Can compile Boolean formula into a decision diagram (OBDD/SDD), such that inference resolves to independent-and and independent-or operations over the decomposed formula.

Top-k pruning [Ré, Davli \& Suciu: ICDE'07; Karp, Luby \& Madras: J-Alg.'89;
Olteanu \& Wen: ICDE'12]

- Can return top-k answers based on lower and upper bounds, even without knowing their exact marginal probabilities.
- Multi-Simulation: run multiple Markov-Chain-Monte-Carlo (MCMC) simulations in parallel.


## Top-k Ranking by Marginal Probabilities



## Bounds for First-Order Formulas

## Theorem 1:

Given a (partially grounded) first-order lineage formula $\Phi$ :

```
\(\Phi\left(\mathrm{Q}_{2}\right)=\mathrm{B} \vee \exists \mathrm{y}\) gradFrom(S,y)
```

- Lower bound $\mathrm{P}_{\text {low }}$ (for all query answers that can be obtained from grounding $\Phi$ ): Substitute $\exists \mathrm{y}$ gradFrom( $\mathrm{S}, \mathrm{y}$ ) with false (or true if negated).

$$
P_{\text {low }}\left(Q_{2}\right)=P(B \vee \text { false })=P(B)=0.6
$$

- Upper bound $P_{\text {up }}$ (for all query answers that can be obtained from grounding $\Phi$ ): Substitute $\exists \mathrm{y}$ gradFrom $(\mathrm{S}, \mathrm{y}$ ) with true (or false if negated).

$$
P_{\text {up }}\left(Q_{2}\right)=P(B \vee \text { true })=P(\text { true })=1.0
$$

Proof: (sketch)
Substitution of a subformula with false reduces the number of models (possible worlds) that satisfy $\Phi$; substitution with true increases them.

## Convergence of Bounds

## Theorem 2:

Let $\Phi_{1}, \ldots, \Phi_{\mathrm{n}}$ be a series of first-order lineage formulas obtained from grounding $\Phi$ via SLD resolution, and let $\varphi$ be the propositional lineage formula of an answer obtained from this grounding procedure.
Then rewriting each $\Phi_{i}$ according to Theorem 1 into $P_{i, \text { low }}$ and $P_{i, \text { up }}$ creates a monotonic series of lower and upper bounds that converges to $\mathrm{P}(\varphi)$.

$$
\begin{aligned}
0=P(\text { false }) \leq P(B \vee \text { false })=0.6 \leq P( & B \vee(C \wedge D))=0.888 \\
& \leq P(B \vee \text { true })=P(\text { true })=1
\end{aligned}
$$

Proof: (sketch, via induction)
Substitution of true with a formula reduces the number of models that satisfy $\Phi$; substitution of false with a formula increases this number.

## Top-k Stopping Condition

[Fagin et al.'01; Balke,Kießling'02; Dylla,Miliaraki,Theobald: ICDE'13]
"Fagin's Algorithm"

- Maintain two disjoint queues:

Top-k sorted by $\mathrm{P}_{\text {low }}$ and Candidates sorted by $\mathrm{P}_{\text {up }}$

- Return the top-k queue at the $t$-th grounding step when:
$\left.P_{\mathrm{i}, \text { low }}\left(\mathrm{Q}_{\mathrm{k}}\right)\right|_{\mathrm{Qk} \in \text { Top-k }}>\left.\mathrm{P}_{\mathrm{i}, \mathrm{up}}\left(\mathrm{Q}_{\mathrm{j}}\right)\right|_{\mathrm{Q}_{\mathrm{j} \in \text { Candidates }}}$



## Temporal-Probabilistic Database

[Wang,Yahya,Theobald: MUD'10; Dylla,Miliaraki,Theobald: PVLDB'13]

Derived Facts

Base Facts
teamMates(Beckham, Recmealthon, $T_{\beta}$ )


Example using the Allen predicate overlaps

## Inference in Temporal-Probabilistic Databases

[Wang,Yahya,Theobald: MUD'10; Dylla,Miliaraki,Theobald: PVLDB'13]


## Inference in Temporal-Probabilistic Databases

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Derived Facts
teamMates(Beckham, teamMates(Beckham,
Ronaldo, $\mathrm{T}_{4}$ ) Zidane, $\mathrm{T}_{5}$ )
teamMates(Ronaldo, Zidane, $\mathrm{T}_{6}$ )

Non-independent
Independent


- Closed and complete representation model (incl. lineage)
- Temporal alignment is linear in the number of input intervals
- Probabilistic inference per interval remains \#P-hard
- Inference requires lineage decompositions, top-k pruning, or Monte Carlo approximations (Luby-Karp for DNF, MCMC-style sampling)


## Ultimate PhD Challenge (II)

Lifted inference with fully integrated relational and probabilistic optimization for arbitrary SQL queries!

- Query answering: graduatedFrom(Surajit, y)

- Boolean queries:
$\exists y$ graduatedFrom(Surajit, y)



## RDF-Centered Research Topics

- Information Extraction
[SIGMOD'09, WebDB'10, PODS'10, WSDM'11, CIKM'12, CLEF/INEX'11''12, LDOW'14, TACL'16]
- Uncertain RDF Data \& Probabilistic Databases
[ICDE'08, VLDB-J'08, SSDBM'10, BTW'11, CIKM'11, ICDE'13, PVLDB'14, VLDB PhD Workshop'15]


## Scalable RDF Indexing \& SPARQL Query

 Processing[SIGMOD'14, SWIM'14, SIGMOD'16]

## RDF \& SPARQL

## RDF Data:

1. Relation(s)
2. Logical facts
3. Directed graphs


Data complexity of core SPARQL: polynomial
Combined data \& query complexity: exponential

RDF
Indexing

## TriAD Architecture




## Stage 2 MPICH2 - Asynchronous Communication Protocol


$\rightarrow$ TriAD follows a very classical master-slave architecture; however with a direct (asynchronous) communication among all slaves at query time.

## Locality-Based Graph Summarization: METIS

$$
k=4
$$

## METIS

- Tools like METIS can efficiently approximate a min- $k$-cut partitioning for graphs with many millions of
 nodes/edges.


## Min-k-Cut

- For a desired amount of $k$ evenly sized partitions, assign each node in the RDF data graph to exactly one partition, such that the number of cut edges among those partitions is minimized.


## Summary Graph



## RDF Summary Graph

- Drop all nodes and edges inside the partitions
- Keep only inter-partition edges
- Introduce self-loop edges for intra-partition edges


## Querying the Summary Graph

SELECT ? C, ?a
WHERE \{
<Barack_Obama> <born> ?c. ?c <located> <USA>.
<Barack_Obama> <won> ?a \}

Global Dictionary:

| Barack_Obama | $\rightarrow P_{1}$ |
| :--- | :--- |
| USA | $\rightarrow P_{1}$ |
| Lady_Gaga | $\rightarrow P_{2}$ |
| Peace_Nobel_Prize | $\rightarrow P_{4}$ |

$\rightarrow P_{1}$
$\rightarrow P_{1}$
_ace_Nobel_Prize
Potential matches!

## Querying the Summary Graph

```
SELECT ?c, ?s
WHERE {
    <Barack_Obama> <born> ?c.
    ?c <located> <USA>.
    <Barack_Obama> <governor> ?s
}
```

Global Dictionary:

| Barack_Obama | $\rightarrow P_{1}$ |
| :--- | :--- |
| USA | $\rightarrow P_{1}$ |
| Lady_Gaga | $\rightarrow P_{2}$ |
| Peace_Nobel_Prize | $\rightarrow P_{4}$ |



- Summary graph guarantees no false negatives (i.e., "missed results"); the subsequent processing of the query against the pruned data graph also ensures no false positives.
- Facilitates join-ahead pruning by skipping over irrelevant partitions.


## Example Query Plan

Cost: $\max (100,10)+5$ Sharding: $\mathbf{R}_{\mathbf{2}}$

| Sharding: $\mathbf{R}_{\mathbf{2}}$ | DIS $\left(\mathbf{R}_{\mathbf{1}}\right)$ | $\mathbf{D I S}\left(\mathbf{R}_{\mathbf{2}}\right)$ |
| :--- | :--- | :--- |
|  | POS | $\operatorname{POS}$ |
| Index: | $\operatorname{POS}$ | $[1]$ |
| Slaves: | $[1,2]$ | $[1]$ |
| Partitions: | $[1,3]$ | 10 |
| Cost: | 100 | 10 |

$\begin{array}{cl}\mathbf{D H J}\left(\mathbf{R}_{1,2,3,4}\right) & \begin{array}{l}\text { Cost: } \max (105,215)+3 \\ \text { Sharding: } \mathbf{R}_{1,2}, \mathbf{R}_{3,4}\end{array}\end{array}$ ? P

$\mathbf{D M J}\left(\mathbf{R}_{3,4}\right)$ Cost: $\max (200,150)+15$ ?p, Sharding: none

DIS( $\left.R_{3}\right) \quad$ DIS( $\left.R_{4}\right)$
PSO PSO
[1,2] [1,2]
[1,2,3]
200
150

```
SELECT ?p, ?c, ?a, ?g
```

SELECT ?p, ?c, ?a, ?g
WHERE {
WHERE {
R1: ?p <born> ?c.
R1: ?p <born> ?c.
R2: ?c <located> <USA>.
R2: ?c <located> <USA>.
R3: ?p <won> ?a.
R3: ?p <won> ?a.
R4: ?p <governor> ?g }

```
R4: ?p <governor> ?g }
```

- A copy of the same query plan is shipped to all slaves:
, DIS operators (leafs) are augmented with locality and pruning information.
- 6 SPO permutations allow the usage of DMJ op's at the first level of joins.


## Distributed \& Multithreaded Query Execution

```
SELECT ?p, ?c, ?a, ?g
WHERE {
R1: ?p <born> ?c.
R2: ?c <located> <USA>.
R3: ?p <won> ?a.
R4: ?p <governor> ?g }
```

- All slaves concurrently and asynchronously process the same query plan, but each over disjoint partitions of the SPO permutation indexes.


Slave 1
Slave 2

## Experiments

TriAD is implemented in C++ using GCC 4.4, Boost-1.5 \& MPICH2. All experiments were run on a proprietary cluster with $32 \times 48$ GB RAM, 2 quad-core XENON CPUs and a 1GBit Ethernet connection.

- LUBM - Lehigh University Benchmark Scale Factor 160: 28 Mio RDF triples $\rightarrow 16$ GB data $\rightarrow 3$ GB index Scale Factor 10240: 1.8 Bio RDF triples $\rightarrow 730$ GB data $\rightarrow 150$ GB index
- BTC - Billion Triples Challenge (2012)

DBpedia/Yago/Freebase: 1.4 Bio RDF triples $\rightarrow 231$ GB data $\rightarrow 130$ GB index

- WSDTS - Waterloo SPARQL Diversity Test Suite

Scale Factor 1000: 109 Mio RDF triples $\rightarrow 15$ GB data $\rightarrow 9.1$ GB index

9 Competitors: RDF-3x (MPII), MonetDB (U-Amsterdam), BitMat (Rensselaer Polytech), TripleBit (U-Huazhong/U-Georgia), Hadoop-RDF-3x (Yale), Apache Hadoop / Spark (UC Berkeley), SHARD (open-source), Trinity.RDF (MSR)

## Benchmark Results

|  | TriAD | $\begin{array}{r} \text { TriAD-SG } \\ (200 \mathrm{~K}) \end{array}$ | Trinity.RDF | SHARD | H-RDF-3X |  | 4store |  | RDF-3X |  | BitMat |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | (cold) | (warm) | (cold) | (warm) | (cold) | (warm) | (cold) | (warm) |
| Q1 | 7,631 | 2,146 | 12,648 | 6.9 E 5 | 2.3 E 6 | 1.7 E 5 | aborted | aborted | 1.9 E 6 | 1.8 E 6 | 17,339 | 11,295 |
| Q2 | 1,663 | 2,025 | 6,018 | 2.1 E 5 | 5.3 E 5 | 4,095 | 1.1 E 5 | 15,113 | 6.3 E 5 | 17,835 | 2.4 E 5 | 1.8 E 5 |
| Q3 | 4,290 | 1,647 | 8,735 | 4.7 E 5 | 2.2E6 | 1.3 E 5 | aborted | aborted | 1.7 E 6 | 1.7 E 6 | 8,429 | 2,679 |
| Q4 | 2.1 | 1.3 | 5 | 3.9 E 5 | 166 | 1 | 1,903 | 12 | 243 | 3 | aborted | aborted |
| Q5 | 0.5 | 0.7 | 4 | 97,545 | 85 | 1 | 2,429 | 12 | 99 | 1 | 472 | 338 |
| Q6 | 69 | 1.4 | 9 | 1.8 E 5 | 5.8 E 5 | 23,440 | 3,572 | 9 | 913 | 287 | 7,796 | 5,377 |
| Q7 | 14,895 | 16,863 | 31,214 | 3.9 E 5 | 2.3 E 6 | 2.1 E 5 | aborted | aborted | 6.5 E 5 | 46,262 | 71,157 | 36,905 |
| Geo- <br> Mean | 24 | 106 | 450 | 3.0 E 5 | 91,378 | 2,406 | - | - | 31,345 | 2,99 |  |  |

LUBM-10240: Query Processing Times in Milliseconds (ms)

|  | \#Results | TriAD | $\begin{array}{r} \text { TriAD-SG } \\ (200 \mathrm{~K}) \end{array}$ | H-RDF-3X (cold) (warm) | $\begin{array}{r} \text { RDF } \\ \text { (cold) } \end{array}$ | F-3X (warm) |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Q1 | 1 | 1.5 | 0.3 | 496 | 297 | 4 |  |  |  |
| Q2 | 1 | 61 |  |  | \#Slaves |  | S1-S7 | F1-F5 | C1-C3 |
| Q3 | 1 | 1 0.6 |  |  | \#Slaves | (Geo.-Mean) | (Geo.-Mean) | (Geo.-Mean) | (Geo.-Mean) |
| Q5 | 5 | 51 |  |  |  |  |  |  |  |
| Q6 | 0 | 0.5 | Tri |  | 1 | 2 | 2 | 94 | 494 |
| Q7 | 0 | 50 |  | ( ${ }^{\text {( }}$-SG(75K) | 1 | 8 | 4 | 35 | 767 |
| Q8 | 292 | 128 | Tri |  | 5 | 2 | 3 | 29 | 270 |
| Geo.- |  |  |  | ARD | 5 | 3.2E5 | 5.8E5 | 7.1 E 5 | 7.7 E 5 |
| Mean | 7.4 |  | RDF-3X (cold) |  | 1 | 10,066 | 167 | 1,749 | 6,610 |
| BTC: Query Prc |  |  |  | RDF-3X (warm) | 1 | 18 | 2 | 41 | 354 |
|  |  |  |  | 1 | 3530 | 10,459 | timeout | timeout |
|  |  |  |  | MonetDB (warm) | 1 | 171 | 744 | timeout | timeout |

WSDTS-1000: Query Processing Times (ms)

## Ultimate PhD Challenge (III)

## - From Map \& Reduce



- over Synchronous Dataflows to Asynchronous Dataflows!



## Summary

## Information Extraction

- Natural-Language

Processing \& Understanding

- Named-Entity Recognition
\& Disambiguation
- Extraction of N-Ary Relations
- Knowledge-Graph

Construction, Integration
\& Maintenance

## Big Data

- Big Data Analytics
- Distributed Graph Engines
- Real-Time Dataflows \& Stream Processing
- Message Passing \& Asynchronous Protocols


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