Einführung in Web und Data ScienceCommunity Analysis

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Today's lecture

- Social Network Analysis
- Anchor text
- Link analysis for ranking
 - PageRank and variants
 - Hyperlink-Induced Topic Search (HITS)



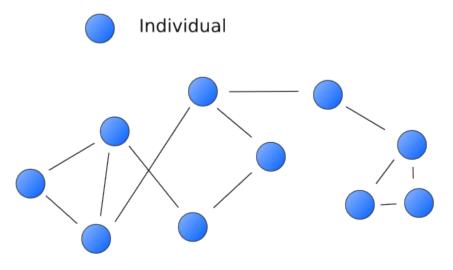
Acknowledgements

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- Thanks also to other lecturers who provided their teaching material on the web



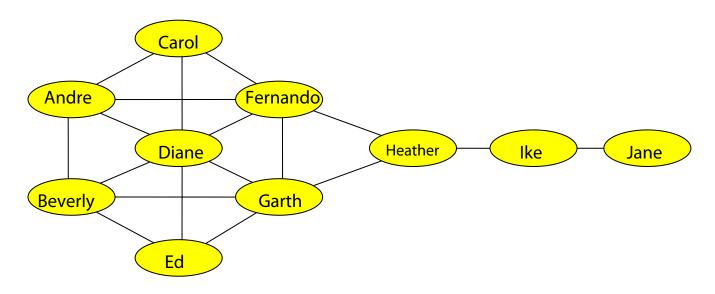
Social Network Analysis (SNA)

- Mapping and measuring of relationships and flows between people, groups, organizations, computers or other information/knowledge processing entities.
- The nodes in the network are the people and groups while the links show relationships or flows between the nodes.





Kite Network



- Who are connecters or hubs in the network?
- Who has control over what flows in the network?
- Who has best visibility of what is happening in the network?
- Who are peripheral players? Are they important?



Measures

1. Degree Centrality:

The number of direct connections a node has. What really matters is where those connections lead to and how they connect the otherwise unconnected.

$$C_D(n_i) = d(n_i)$$

$$C_D'(n_i) = \frac{d(n_i)}{g-1}$$

2. Betweenness Centrality:

A node with high betweenness has great influence over what flows in the network indicating important links and single points of failure.

Rescaling by dividing through by the number of pairs of nodes not including n_i

$$C_B(n_i) = \sum_{j < k} g_{jk}(n_i) / g_{jk}$$
 $C'_B(n_i) = \frac{C_B(n_i)}{(g-1)(g-2)/2}$

3. Closeness Centrality:

The measure of closeness of a node to everyone else.

Determined by the sum of the length of the <u>shortest paths</u> between the node and all other nodes in the graph.

Undirected graph

$$C_{C}(n_{i}) = \left[\sum_{j=1}^{g} d(n_{i}, n_{j})\right]^{-1}$$

$$C'_{C}(n_{i}) = \frac{g-1}{\sum_{i=1}^{g} d(n_{i}, n_{j})} = (g-1)C_{C}(n_{i})$$

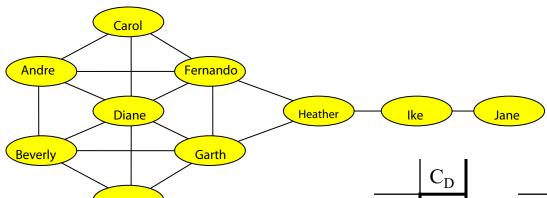


Legend

- g = size of graph (number of nodes)
- d(.) = (in)degree
- g_{ik} = number of minimal paths between nodes j and k
- g_{jk}(n) = number of minimal paths between nodes j and k
 that contain n
- (g-1)(g-2)/2 = number of potential paths $\Sigma_{x=1}^{u} x = (u+1)u/2 \text{ für } u=(g-2)$
- d(.,.)= distance between two nodes
- Scaling with (g 1)(g 2): For every node n except n_i pair the node with all other nodes except n and n_i



Example: Kite-Network



$$C_B(n_i) = \sum_{j < k} g_{jk}(n_i) / g_{jk}$$

$$C_{C}(n_{i}) = \left[\sum_{j=1}^{g} d(n_{i}, n_{j})\right]^{-1}$$

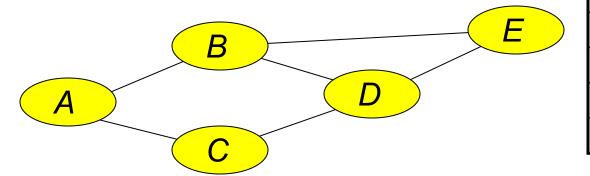
$$C_D(n_i) = d(n_i)$$

	C_{D}
С	3
A	4
F	5
D	6
В	4
G	5
Е	3
Н	3
Ι	2
J	1

	C	A	F	D	В	G	Е	Н	I	J
С	0	1	1	1	0	0	0	0	0	0
A	1	0	1	1	1	0	0	0	0	0
F	1	1	0	1	0	1	0	1	0	0
D	1	1	1	0	1	1	1	0	0	0
В	0	1	0	1	0	1	1	0	0	0
G	0	0	1	1	1	0	1	1	0	0
Е	0	0	0	1	1	1	0	0	0	0
Н	0	0	1	0	0	1	0	0	1	0
I	0	0	0	0	0	0	0	1	0	1
J	0	0	0	0	0	0	0	0	1	0

Example

$$C_B(n_i) = \sum_{j < k} g_{jk}(n_i) / g_{jk}$$
 $C_C(n_i) = \left[\sum_{j=1}^g d(n_i, n_j)\right]^{-1}$ $C_D(n_i) = d(n_i)$



	A	В	С	D	Е
A	0	1	1	0	0
В	1	0	0	1	1
C	1	0	0	1	0
D	0	1	1	0	1
Е	0	1	0	1	0

Adjacency

	C_{B}	C_{C}	C_{D}
A	1	1/6	2
В	3	1/5	3
С	1	1/6	2
D	3	1/5	3
Е	0	1/6	2

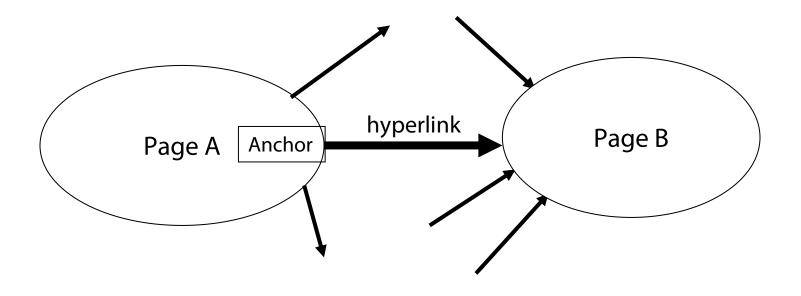
	A	В	С	D	Е
A	0	1	1	2	2
В	1	0	2	1	1
С	1	2	0	1	2
D	2	1	1	0	1
Е	2	1	2	1	0

	A	В	С	D	Е
A	0	A	A	BC	В
В	В	0	AD	В	В
С	С	AD	0	С	D
D	BC	D	D	0	D
Е	В	Е	D	Е	0

Distance

Paths

The Web as a Directed Graph



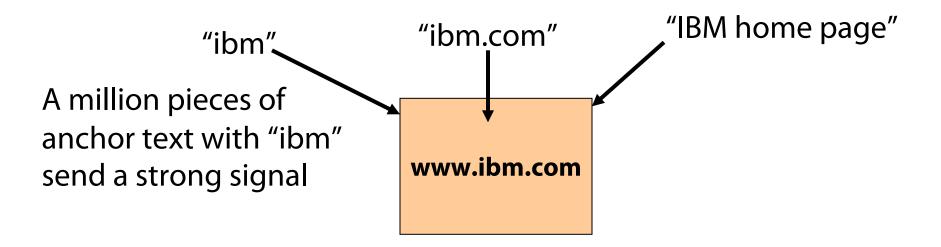
Assumption 1: A hyperlink between pages denotes author perceived relevance (quality signal)

Assumption 2: The anchor of the hyperlink describes the target page (textual context)



Anchor Text

- For IBM how to distinguish between:
 - IBM's home page (mostly graphical)
 - IBM's copyright page (high term freq. for 'ibm')
 - Rival's spam page (arbitrarily high term freq.)

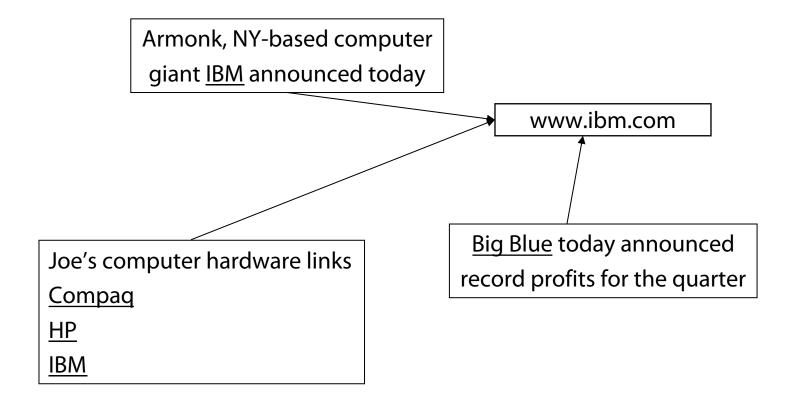


Oliver A. McBryan. GENVL and WWWW: Tools for Taming the Web. Research explained at First International Conference on the World Wide Web. CERN, Geneva (Switzerland), May 25-26-27 **1994** (WWWW=World Wide Web Worm, first serach engine for the web)



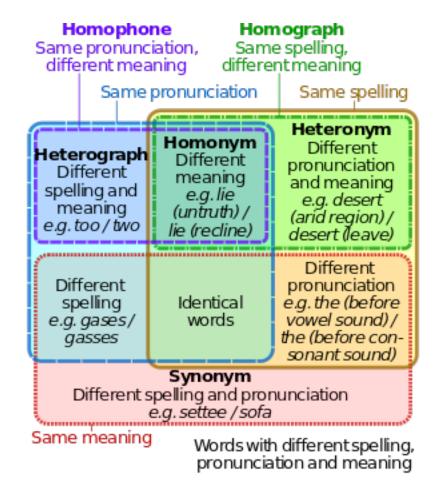
Indexing anchor text

 When indexing a document D, include anchor text from links pointing to D.





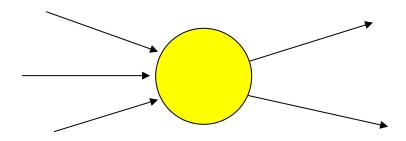
The Web as a Resource for NLP





The Web as a Resource for Ranking

- First generation: using link counts as simple measures of popularity.
- Two basic suggestions:
 - Undirected popularity:
 - Each page gets a score = the number of in-links plus the number of out-links (3+2=5).
 - Directed popularity:
 - Score of a page = number of its in-links (3).





Query processing

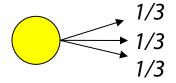
- First retrieve all pages matching the text query (say *venture capital*).
- Order these by their link popularity (either variant on the previous page).

How to organize for "Search Engine Optimization"?



PageRank scoring

- Imagine a browser doing a random walk on web pages:
 - Start at a random page

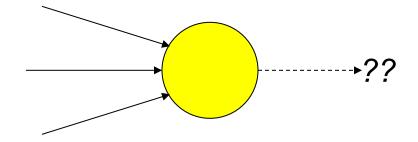


- At each step, go out of the current page along one of the links on that page, equiprobably
- Each page has a long-term visit rate use this as the page's score



Not quite enough

- The web is full of dead-ends.
 - Random walk can get stuck in dead-ends.
 - Makes no sense to talk about long-term visit rates.



Teleporting / damping

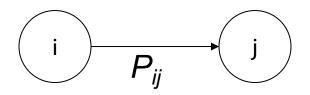
- At a dead end, jump to a random web page.
- At any non-dead end, with probability 10%, jump to a random web page.
 - With remaining probability (90%), go out on a random link.
 - 10% a parameter.
- There is a long-term rate at which any page is visited.
 - How do we compute this visit rate?



Markov chains

- A Markov chain consists of n states, plus an $n \times n$ transition matrix **P**.
- At each step, we are in exactly one of the states.
- For $1 \le i,j \le n$, the matrix entry P_{ij} tells us the relative frequency of j being the next state, given we are currently in state i.

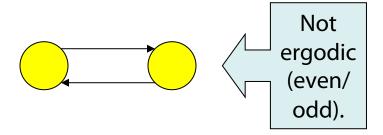
$$\sum_{j=1}^n P_{ij} = 1.$$





Ergodic Markov chains

- A Markov chain is <u>ergodic</u> if
 - you have a path from any state to any other (reducibility)
 - returns to states occur at irregular times (aperiodicity)
 - For any start state, after a finite transient time T_0 , the probability of being in any state at a fixed time $T>T_0$ is nonzero. (positive recurrence)





Ergodic Markov chains

- For any ergodic Markov chain, there is a unique long-term visit rate for each state.
 - "Steady-state" distribution.
- Over a long time-period, we visit each state in proportion to this rate.
- It doesn't matter where we start.



State vectors

- A (row) vector (state vector) $\mathbf{x} = (x_1, ..., x_n)$ tells us where the walk is at any point.
- E.g., (000...1...000) means we're in state i.

1 i n

More generally, the vector $\mathbf{x} = (x_1, ..., x_n)$ means the walk is in state i with relative frequency x_i .

$$\sum_{i=1}^n x_i = 1.$$

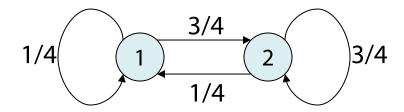
Change in state vector

- If the state vector is $\mathbf{x} = (x_1, \dots x_n)$ at this step, what is it at the next step?
- Recall that row i of the transition matrix P tells us where we go next from state i
- So from x, our next state is distributed as xP.



Steady state example

- The steady state looks like a vector of probabilities $\mathbf{a} = (a_1, \dots a_n)$:
 - $-a_i$ is the relative frequency that we are in state i.



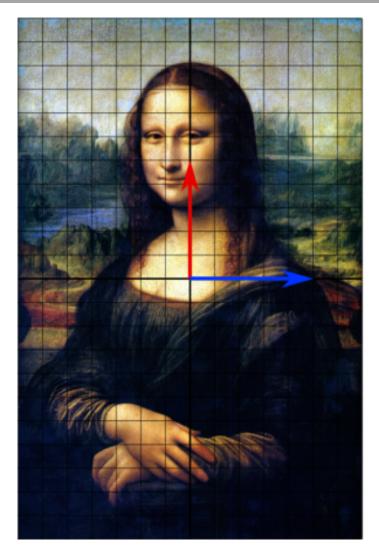
For this example, $a_1=1/4$ and $a_2=3/4$.

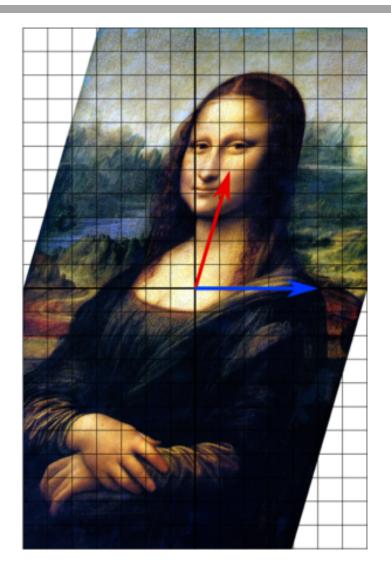
How do we compute this vector?

- Let $\mathbf{a} = (a_1, \dots a_n)$ denote the row vector of steady-state rates.
- If we our current position is described by **a**, then the next step is described as **aP**.
- But a is the steady state, so a=aP.
- Solving this matrix equation gives us a.
 - So a is the (left) eigenvector for P.
 - (Corresponds to the "principal" eigenvector of **P** with the largest eigenvalue)
 - Transition matrices always have largest eigenvalue 1.



Eigenvectors and Eigenvalues $Mx = \lambda x$







One way of computing a

- Recall, regardless of where we start, we eventually reach the steady state a.
- Start with any distribution (say $\mathbf{x} = (10...0)$).
- After one step, we're at xP;
- after two steps at \mathbf{xP}^2 , then \mathbf{xP}^3 and so on.
- "Eventually" means for "large" k, $\mathbf{x}\mathbf{P}^k = \mathbf{a}$.
- Algorithm: multiply x by increasing powers of P until the product looks stable.



PageRank Summary

- Preprocessing:
 - Given graph of links, build matrix P
 - From it compute a
 - The entry a_i is a number between 0 and 1: the pagerank of page i.
- Query processing:
 - Retrieve pages meeting query
 - Rank them by their pagerank
 - Order is query-independent
- A variant of PageRank is used in Google, but also many other clever heuristics



PageRank: Issues and Variants

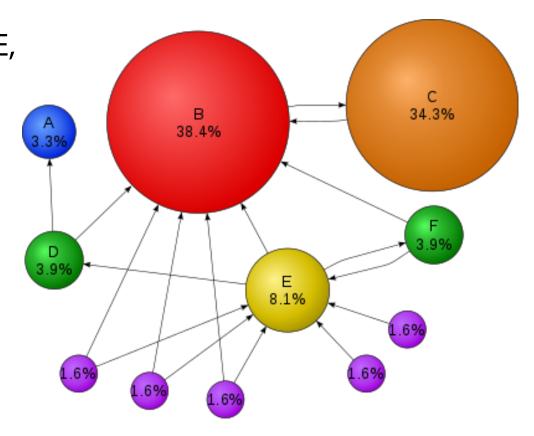
- How realistic is the random surfer model?
 - What if we modeled the back button?
 - Surfer behavior sharply skewed towards short paths
 - Search engines, bookmarks & directories make jumps non-random
- Biased Surfer Models
 - Weight edge traversal probabilities based on match with topic/query (non-uniform edge selection)
 - Bias jumps to pages on topic (e.g., based on personal bookmarks & categories of interest)



Google PageRank

 Links are also weighted according to the importance of the source node

Page C has a higher
 PageRank than Page E,
 even though there
 are fewer links to C;
 the one link to C
 comes from an
 important page
 and hence is
 of high value.





Hyperlink-Induced Topic Search (HITS)

- In response to a query, instead of an ordered list of pages each matching the query, find two sets of inter-related pages:
 - Hub pages are good lists of links on a subject
 - e.g., "Bob's list of cancer-related links."
 - Authority pages occur recurrently on good hubs for the subject
- Best suited for "broad topic" queries rather than for page-finding queries

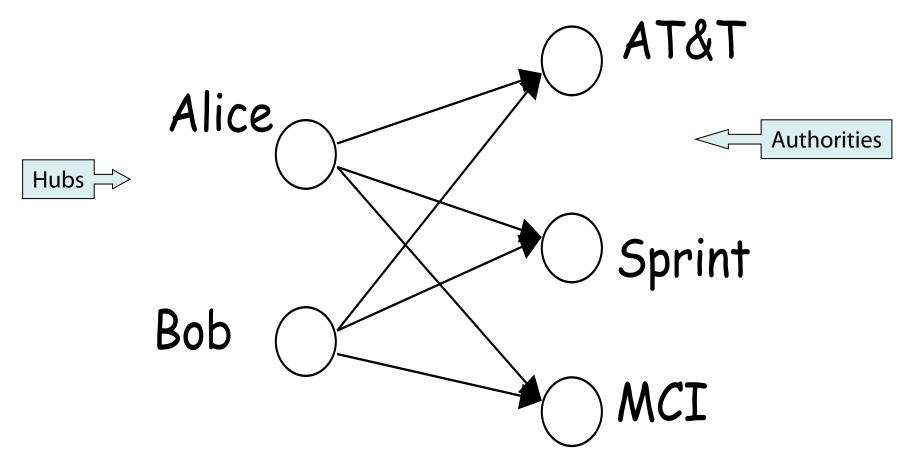


Hubs and Authorities

- Thus, a good hub page for a topic points to many authoritative pages for that topic
- A good authority page for a topic is pointed to by many good hubs for that topic
- Circular definition will turn this into an iterative computation



The hope



Long distance telephone companies



High-level scheme

- Extract from the web a <u>base set</u> of pages that could be good hubs or authorities
- From these, identify a small set of top hub and authority pages;
 - →iterative algorithm

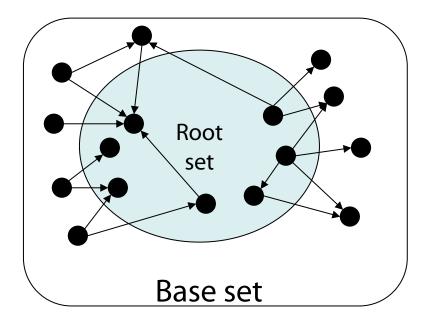


Base set

- Given text query (say "browser"), use a text index to get all pages containing "browser"
 - Call this the root set of pages
- Add in any page that either
 - points to a page in the root set, or
 - is pointed to by a page in the root set
- Call this the base set



Visualization





Assembling the base set

- Root set typically 200-1000 nodes
- Base set may have up to 5000 nodes
- How do you find the base set nodes?
 - Follow out-links by parsing root set pages
 - Get in-links (and out-links) from a connectivity server
 - Actually, suffices to text-index strings of the form
 href="URL" to get in-links to URL



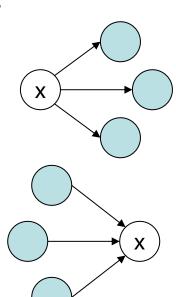
Distilling hubs and authorities

Compute, for each page x in the base set, a hub score h(x) and an authority score a(x)

- Initialize: for all x, $h(x) \leftarrow 1$; $a(x) \leftarrow 1$;
- Iteratively update all h(x), a(x);

$$h(x) \leftarrow \sum_{x \mapsto y} a(y)$$

$$a(x) \leftarrow \sum_{y \mapsto x} h(y)$$



- After iterations output pages with
 - highest h() scores as top hubs
 - highest a() scores as top authorities

Scaling

- To prevent the h() and a() values from getting too big, can scale down after each iteration
- Scaling factor doesn't really matter:
 - we only care about the relative values of the scores



How many iterations?

- Claim: relative values of scores will converge after a few iterations:
 - In fact, suitably scaled, h() and a() scores settle into a steady state!
- We only require the relative orders of the h() and a()
 scores not their absolute values
- In practice, ~5 iterations get you close to stability



Tiefes Verstehen

- Bestimmung von bedingten Unabhangigkeiten zwischen Attributwerten von Objekten in einem sozialen Netz
- Bestimmung von Clustern

•



Data Models vs. Algorithmic Models

