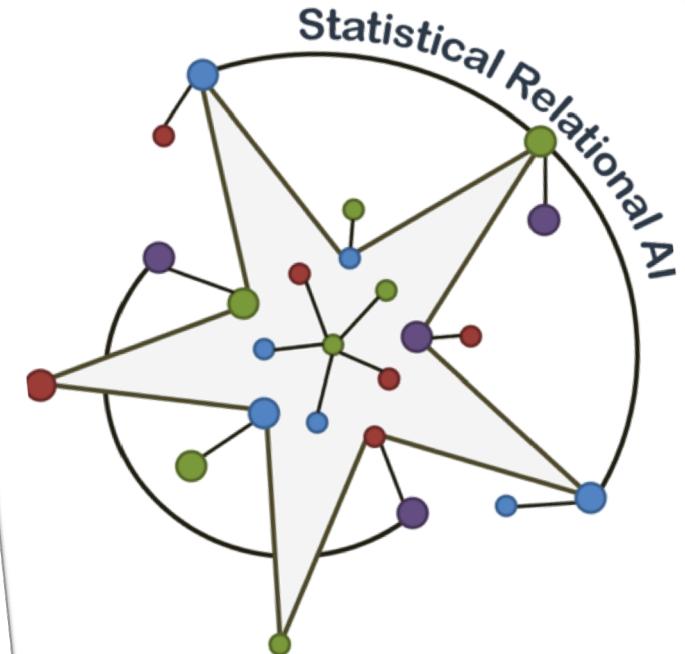


Probabilistic Relational Modeling

Statistical Relational AI

Tutorial at KI-2018



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Agenda: Probabilistic Relational Modeling

- Application
 - Information retrieval (IR)
 - Probabilistic Datalog
- Probabilistic relational logics
 - Overview
 - Semantics
 - Inference problems
- Scalability issues
 - Proposed solutions



Application

Probabilistic Datalog for information retrieval [Fuhr 95]:

```
0.7 term(d1, ir).  
0.8 term(d1, db).  
0.5 link(d2, d1).  
about(D, T) :- term(D, T).  
about(D, T) :- link(D, D1), about(D1, T).
```

Query/Answer:

```
: - term(X, ir) & term(X, db).  
X = 0.56 d1
```

Application: Probabilistic IR

Probabilistic Datalog:

```
0.7 term(d1, ir).  
0.8 term(d1, db).  
0.5 link(d2, d1).  
about(D, T) :- term(D, T).  
about(D, T) :- link(D, D1), about(D1, T).
```

Query/Answer:

```
q(X) :- term(X, ir).  
q(X) :- term(X, db).  
:- q(X)  
X = 0.94 d1
```

Application: Probabilistic IR

Probabilistic Datalog:

```
0.7 term(d1, ir).  
0.8 term(d1, db).  
0.5 link(d2, d1).  
about(D, T) :- term(D, T).  
about(D, T) :- link(D, D1), about(D1, T).
```

Query/Answer:

```
: - about(X, db).  
X = 0.8 d1;  
X = 0.4 d2
```

Application: Probabilistic IR

Probabilistic Datalog:

```
0.7 term(d1, ir).  
0.8 term(d1, db).  
0.5 link(d2, d1).  
about(D, T) :- term(D, T).  
about(D, T) :- link(D, D1), about(D1, T).
```

Query/Answer:

```
: - about(X, db) & about(X, ir).  
X = 0.56 d1;  
X = 0.28 d2 # NOT naively 0.14 = 0.8*0.5*0.7*0.5
```

Solving Inference Problems

- QA requires **proper probabilistic reasoning**
- **Scalability** issues
 - Grounding and propositional reasoning?
 - In this tutorial the focus is on **lifted reasoning** in the sense of [Poole 2003]
 - Lifted exact reasoning
 - Lifted approximations
- Need an overview of the field:
Consider related approaches first

D. Poole. "First-order Probabilistic Inference." In: IJCAI-03
Proceedings of the 18th International Joint Conference on
Artificial Intelligence. 2003

Application: Probabilistic IR

Uncertain Datalog rules: Semantics?

0.7 term(d1, ir) .

0.8 term(d1, db) .

0.5 link(d2, d1) .

0.9 about(D, T) :- term(D, T) .

0.7 about(D, T) :- link(D, D1), about(D1, T) .

Application: Probabilistic IR

Uncertain Datalog rules: Semantics?

0.7 term(d1, ir).

0.8 term(d1, db).

0.5 link(d2, d1).

0.9 temp1.

0.7 temp2.

about(D, T) :- term(D, T), temp1.

about(D, T) :- link(D, D1), about(D1, T), temp2.

Probabilistic Datalog: QA

- Derivation of lineage formula
with Boolean variables corresponding to used facts

Thomas Rölleke; Norbert Fuhr, Information Retrieval with Probabilistic Datalog. In: Logic and Uncertainty in Information Retrieval: Advanced models for the representation and retrieval of information, **1998**.

O. Benjelloun, A. Das Sarma, A. Halevy, M. Theobald, and J. Widom. Databases with Uncertainty and Lineage. VLDB Journal, 17(2):243-264, March **2008**.

- Probabilistic relational algebra

N. Fuhr; T. Rölleke, A Probabilistic Relational Algebra for the Integration of Information Retrieval and Database Systems. ACM Transactions on Information Systems 14(1), **1997**.

- Ranking / top-k QA

N. Fuhr. 2008. A probability ranking principle for interactive information retrieval. Inf. Retr. 11, 3, 251-265, **2008**.

Probabilistic Relational Logics: Semantics

- Distribution semantics (aka grounding or Herbrand semantics) [Sato 95]
Completely define discrete joint distribution by "factorization"
Logical atoms treated as random variables
 - Probabilistic extensions to Datalog [Schmidt et al. 90, Dantsin 91, Ng & Subramanian 93, Poole et al. 93, Fuhr 95, Rölleke & Fuhr 97 and later]
 - Primula [Jaeger 95 and later]
 - BLP, ProbLog [De Raedt, Kersting et al. 07 and later]
 - Probabilistic Relational Models (PRMs) [Poole 03 and later]
 - Markov Logic Networks (MLNs) [Domingos et al. 06]
- Probabilistic Soft Logic (PSL) [Kimmig, Bach, Getoor et al. 12]
Define density function using log-linear model
- Maximum entropy semantics [Kern-Isberner, Beierle, Finthammer, Thimm 10, 12]
Partial specification of discrete joint with "uniform completion"

Additional Bibliography for Prev Slide

Schmidt, H.; Steger, N.; Güntzer, U.; Kiessling, W.; Azone, R.; Bayer, R.,
Combining Deduction by Certainty with the Power of Magic. In: Deductive and
Object-Oriented Databases, pages 103–122. Elsevier Science Publishers,
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Conference on Logic Programming, volume 592 of LNCS, pages 152–164.
Springer, **1991**.

Ng, R.; Subrahmanian, V. S., A Semantical Framework for Supporting
Subjective and Conditional Probabilities in Deductive Databases. Journal of
Automated Reasoning 10, pages 191–235, **1993**.

Poole, D., Logic Programming, Abduction and Probability. New Generation
Computing 11(3), pages 377–400, **1993**.

Inference Problems w/ and w/o Evidence

- Static case
 - Projection (margins),
 - Most-probable explanation (MPE)
 - Maximum a posterior (MAP)
 - Query answering (QA): compute bindings
- Dynamic case
 - Filtering (current state)
 - Prediction (future states)
 - Hindsight (previous states)
 - MPE, MAP (temporal sequence)

ProbLog

```
% Intensional probabilistic facts:
```

```
0.6::heads(C) :- coin(C).
```

```
% Background information:
```

```
coin(c1).
```

```
coin(c2).
```

```
coin(c3).
```

```
coin(c4).
```

```
% Rules:
```

```
someHeads :- heads(_).
```

```
% Queries:
```

```
query(someHeads).
```

```
0.9744
```

- Compute marginal probabilities of any number of ground atoms in the presence of evidence
- Learn the parameters of the ProbLog program from partial interpretations
- Sample from a ProbLog program
 - Generate random structures (use case: [Goodman & Tenenbaum 16])
- Solve decision theoretic problems:
 - Decision facts and utility statements

ProbLog: A probabilistic Prolog and its application in link discovery, L. De Raedt, A. Kimmig, and H. Toivonen, Proceedings of the 20th International Joint Conference on Artificial Intelligence (IJCAI-07), Hyderabad, India, pages 2462-2467, 2007

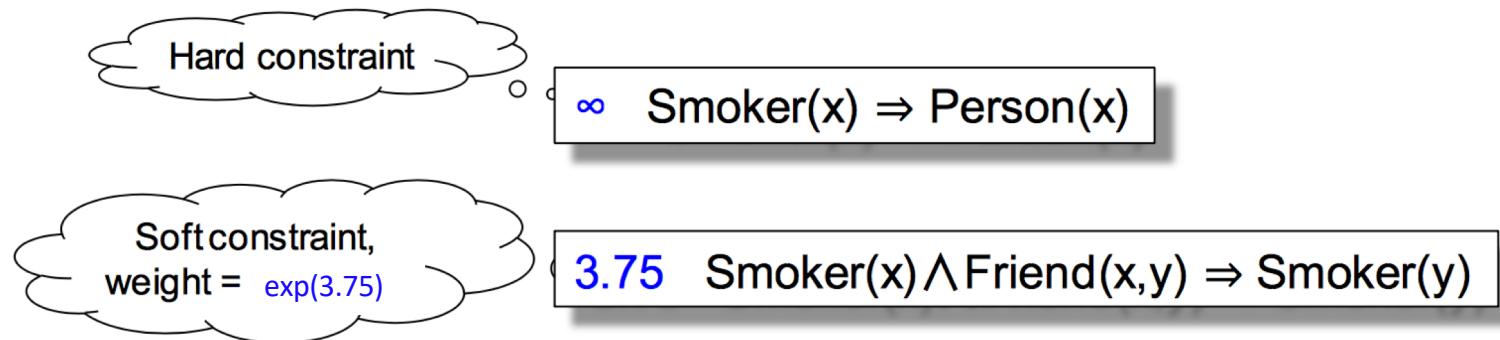
K. Kersting and L. De Raedt, Bayesian logic programming: Theory and Tool. In L. Getoor and B. Taskar, editors, An Introduction to Statistical Relational Learning. MIT Press, 2007

Daan Fierens, Guy Van den Broeck, Joris Renkens, Dimitar Shterionov, Bernd Gutmann, Ingo Thon, Gerda Janssens, and Luc De Raedt. Inference and learning in probabilistic logic programs using weighted Boolean formulas, In: Theory and Practice of Logic Programming, 2015

N. D. Goodman, J. B. Tenenbaum, and The ProbMods Contributors. Probabilistic Models of Cognition (2nd ed.), 2016.
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Markov Logic Networks (MLNs)

- Weighted formulas for modeling constraints [Richardson & Domingos 06]



An MLN is a set of constraints $(w, \Gamma(x))$, where w =weight, $\Gamma(x)$ =FO formula

Weight of a world = product of $\exp(w)$, for all MLN rules $(w, \Gamma(x))$ and groundings $\Gamma(a)$ that hold in that world

Probability of a world = Weight / Z

Z = sum of weights of all worlds (no longer a simple expression!)

Why exp? Semantics of MLNs

- Let D be a set of constants ...
- ... and $\omega \in \{0,1\}^m$ a **world** with m atoms w.r.t. D
- $weight(\omega) = \prod_{\{(w, \Gamma(x)) \in MLN \mid \exists a \in D^n : \omega \models \Gamma(a)\}} \exp(w)$
- $\ln(weight(\omega)) = \sum_{\{(w, \Gamma(x)) \in MLN \mid \exists a \in D^n : \omega \models \Gamma(a)\}} w$
 - Sum allows for component-wise optimization
(e.g., during weight learning)
- $Z = \sum_{\omega \in \{0,1\}^m} \ln(weight(\omega))$
- $P(\omega) = \ln(weight(\omega)) / Z$

Factorization

Log-Linear

MLN Query answering semantics $P(Smoker(john)) =$

$$\sum_{\omega \in \{0,1\}^m \wedge \omega \models Smoker(John)} P(\omega)$$

Maximum Entropy Principle

- Given:
 - States $s = s_1, s_2, \dots, s_n$
 - Density $p_s = (p_1, p_2, \dots, p_n)$
- Maximum Entropy Principle:
 - W/o further information, select p_s ,
s.t. extropy is maximized
$$-\sum_s p_s(s) \log p_s(s) = -p_s \log p_s$$
 - w.r.t. k constraints (expected values of features)

$$\sum_s p_s(s) f_i(s) = D_i \quad i \in \{1..k\}$$

Maximum Entropy Principle

- Maximize Lagrange functional for determining p_s

$$L = -p_s \log p_s - \sum_i \lambda_i (\sum_s p_s(s) f_i(s) - D_i) - \mu (\sum_s p_s(s) - 1)$$

- Partial derivatives of L w.r.t. $p_s \rightarrow$ roots:

$$p_s(s) = \frac{\exp\left(-\sum_i \lambda_i f_i(s)\right)}{Z}$$

where Z is for normalization
(Boltzmann-Gibbs distribution)

- "Global" modeling: additions/changes to constraints/rules influence the whole joint probability distribution

Maximum-Entropy Semantics for PRMs

- Probabilistic Conditionals [Kern-Isberner et al 10, 12]

$r_1 : (likes(X, Y) | el(X) \wedge ke(Y))[0.6]$

$r_2 : (likes(X, fred) | el(X) \wedge ke(fred))[0.4]$

$r_3 : (likes(clyde, fred) | el(clyde) \wedge ke(fred))[0.7]$

$el = elephant, ke = keeper$

- Averaging semantics
- Aggregating semantics
- Allows for "local modeling" → transfer learning made easier

G. Kern-Isberner and M. Thimm. "Novel Semantical Approaches to Relational Probabilistic Conditionals." In: Proc. KR'10, pp. 382–392, 2010.

G. Kern-Isberner, C. Beierle, M. Finthammer, and M. Thimm. "Comparing and Evaluating Approaches to Probabilistic Reasoning: Theory, Implementation, and Applications." In: Transactions on Large-Scale Data- and Knowledge-Centered Systems VI. LNCS 7600. Springer , pp. 31–75, 2012 .

M. Finthammer, "Concepts and Algorithms for Computing Maximum Entropy Distributions for Knowledge Bases with Relational Probabilistic Conditionals." IOS Press, 2017.

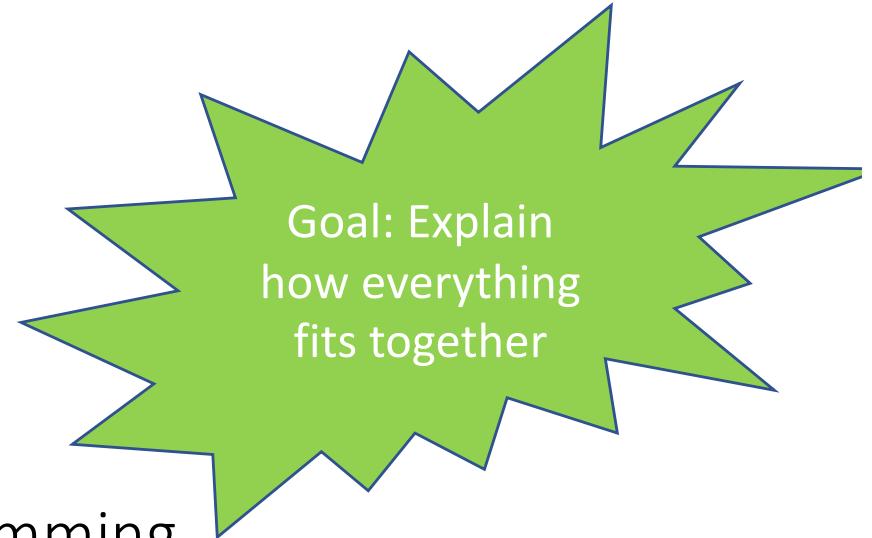
Factor graphs

- Unifying representation for specifying discrete distributions with a factored representation
 - Potentials (weights) rather than probabilities
- Also used in engineering community for defining densities w.r.t. continuous domains
[Loeliger et al. 07]

Research on Scalability

Proposed solutions

- Limited expressivity:
 - Probabilistic databases
- Knowledge Compilation
 - Transformation to linear programming
 - Weighted first-order model counting
- Approximation
 - Grounding + belief propagation (TensorLog)

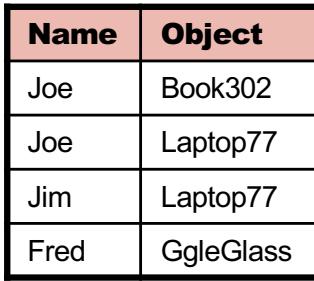
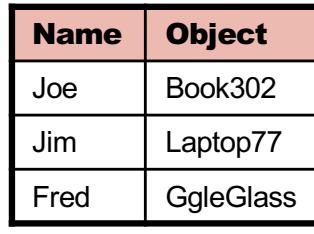
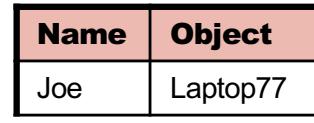
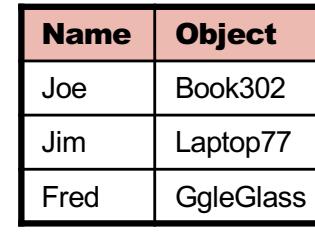
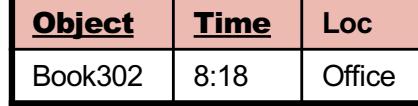
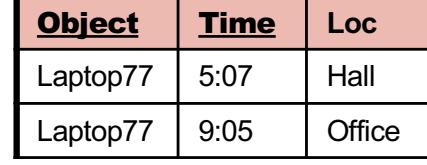
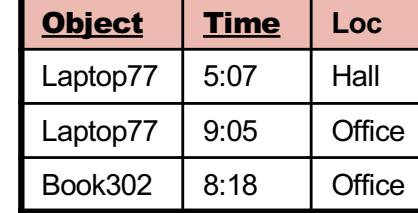
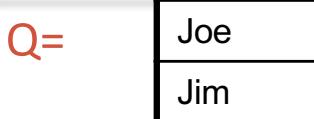


Probabilistic Databases

$$P(\text{Joe}) = 1.0$$

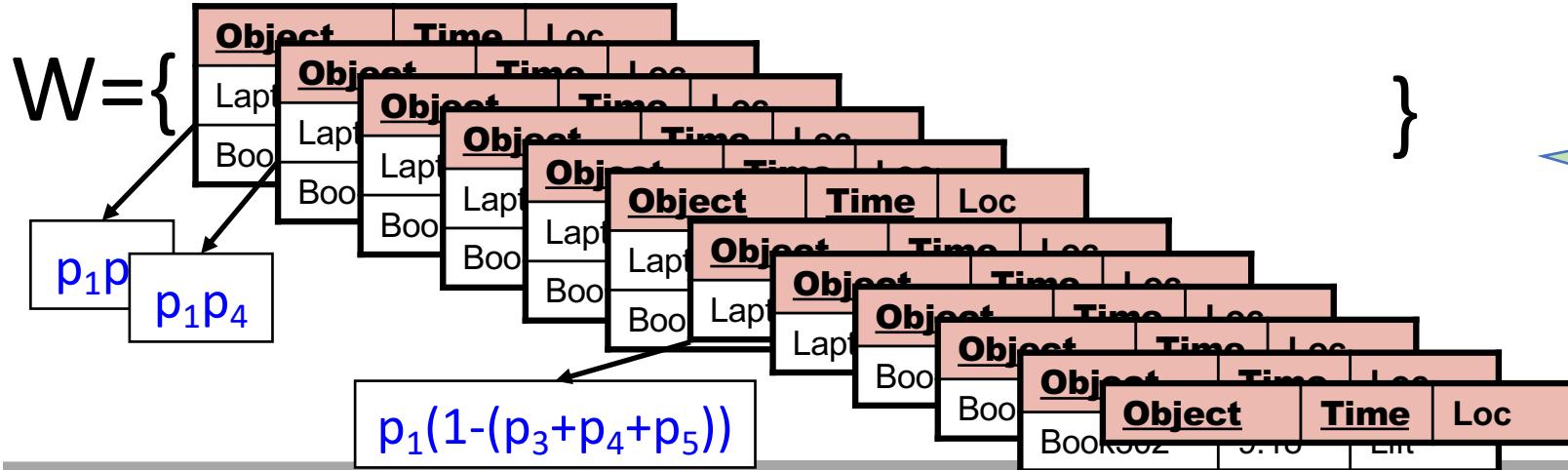
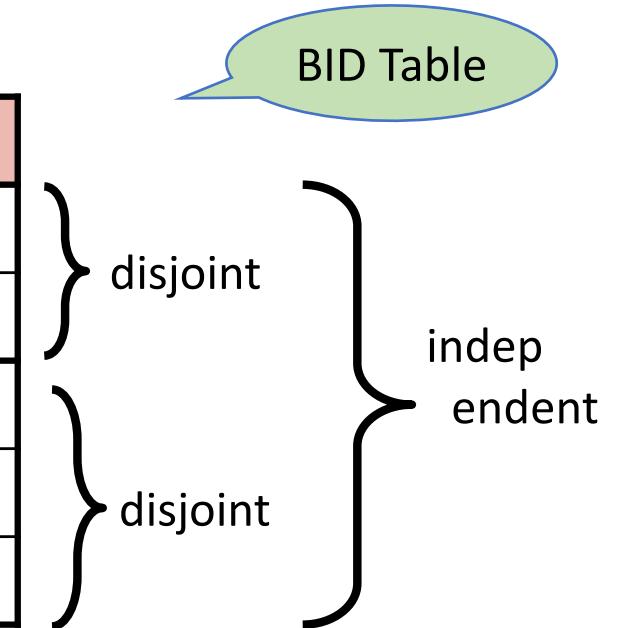
$$P(\text{Jim}) = 0.4$$

$$Q(z) = \text{Owner}(z,x), \\ \text{Location}(x,t,\text{'Office'})$$

W_1	W_2	W_3	W_4
Owner  0.3	Owner  0.4	Owner  0.2	Owner  0.1
Location 	Location 	Location 	Location 
$Q =$ 	$Q =$ 	$Q =$ 	$Q =$ 

BID Tables

Object	Time	Loc	P
Laptop77	9:07	Rm444	p_1
Laptop77	9:07	Hall	p_2
Book302	9:18	Office	p_3
Book302	9:18	Rm444	p_4
Book302	9:18	Lift	p_5



QA: Example

Transformation to SQL
is possible

```
SELECT DISTINCT 'true'  
FROM R, S  
WHERE R.x = S.x
```

$Q() = R(x), S(x,y)$

$$P(Q) = 1 - \{1 - p1 * [1 - (1-q1) * (1-q2)]\} * \\ \{1 - p2 * [1 - (1-q3) * (1-q4) * (1-q5)]\}$$

Determine $P(Q)$ in PTIME
w.r.t. size of database

R

x	P
a1	p1
a2	p2
a3	p3

S

x	y	P
a1	b1	q1
a1	b2	q2
a2	b3	q3
a2	b4	q4
a2	b5	q5

Problem: Some Queries don't Scale

- Dichotomy P vs. #P [Suciu et al. 11]
- Important research area: **Classes of liftable queries**
 - Transform queries to SQL based on rule set [Suciu et al. 11]
 - If transformation possible: SQL might do the job
 - Non-recursive Datalog: Expand queries, then transform
 - Further query classes that allow for lifting [Kazemi et al. 17]
 - Not all queries are liftable
 - With probabilistic databases, transformed queries tend to be large and complex

Transformation to SQL does not work for the IR example from the beginning

No guarantee that SQL engines can handle them

D. Suciu, D. Olteanu, R. Christopher, and C. Koch. Probabilistic Databases. 1st. Morgan & Claypool Publishers, 2011.

S. M. Kazemi, A. Kimmig, G. Van den Broeck, and D. Poole. "Domain Recursion for Lifted Inference with Existential Quantifiers." In: Seventh International Workshop on Statistical Relational AI (StarAI). Aug. 2017

Probabilistic Relational Logic

- First-order logic formulas for **expressivity**
- Knowledge compilation for **scalability**
 - Compilation to linear programming
 - Probabilistic Soft Logic [Kimmig, Bach, Getoor et al. 12]
 - Probabilistic Doxastic Temporal Logic [Martiny & Möller 16]
 - Weighted first-order model counting (WFOMC)
[Van den Broeck, G., Taghipour, N., Meert, W., Davis, J., & De Raedt, L.]

Kimmig, A., Bach, S. H., Broecheler, M., Huang, B. & Getoor, L. A Short Introduction to Probabilistic Soft Logic. NIPS Workshop on Probabilistic Programming: Foundations and Applications, 2012.

Karsten Martiny, Ralf Möller: PDT Logic: A Probabilistic Doxastic Temporal Logic for Reasoning about Beliefs in Multi-agent Systems
In: J. Artif. Intell. Res. (JAIR), Vol.57, p.39-112, 2016.

Van den Broeck, G., Taghipour, N., Meert, W., Davis, J., & De Raedt, L.,
Lifted probabilistic inference by first-order knowledge compilation.
In Proc.IJCAI-11, pp. 2178-2185, 2011.

Probabilistic Soft Logic (PSL): Example

- First Order Logic weighted rules

$$0.3 : \text{friend}(B, A) \wedge \text{votesFor}(A, P) \rightarrow \text{votesFor}(B, P)$$
$$0.8 : \text{spouse}(B, A) \wedge \text{votesFor}(A, P) \rightarrow \text{votesFor}(B, P)$$

- Evidence

$$\text{I}(\text{friend}(\text{John}, \text{Alex})) = 1 \quad \text{I}(\text{spouse}(\text{John}, \text{Mary})) = 1$$
$$\text{I}(\text{votesFor}(\text{Alex}, \text{Romney})) = 1 \quad \text{I}(\text{votesFor}(\text{Mary}, \text{Obama})) = 1$$

- Inference

- $\text{I}(\text{votesFor}(\text{John}, \text{Obama})) = 1$
- $\text{I}(\text{votesFor}(\text{John}, \text{Romney})) = 0$

PSL's Interpretation of Logical Connectives

- Continuous truth assignment
Łukasiewicz relaxation of AND, OR, NOT
 - $I(\ell_1 \wedge \ell_2) = \max \{0, I(\ell_1) + I(\ell_2) - 1\}$
 - $I(\ell_1 \vee \ell_2) = \min \{1, I(\ell_1) + I(\ell_2)\}$
 - $I(\neg \ell_1) = 1 - I(\ell_1)$
- Distance to satisfaction d
 - Implication: $\ell_1 \rightarrow \ell_2$ is satisfied iff $I(\ell_1) \leq I(\ell_2)$
 - $d = \max \{0, I(\ell_1) - I(\ell_2)\}$
- Example
 - $I(\ell_1) = 0.3, I(\ell_2) = 0.9 \Rightarrow d = 0$
 - $I(\ell_1) = 0.9, I(\ell_2) = 0.3 \Rightarrow d = 0.6$

PSL Probability Distribution

- Density function:

$$f(I) = \frac{1}{Z} \exp\left[- \sum_{r \in R} \lambda_r (d_r(I))^\perp\right]$$

a possible continuous truth assignment

Normalization constant

For all rules

Weight of formula r

Distance to satisfaction of rule r

One of the features discussed above

Weighted First-Order Model Counting (WFOMC)

Model = Satisfying assignment of a propositional formula Δ

$$\Delta = \forall d (\text{Rain}(d) \Rightarrow \text{Cloudy}(d))$$

Days = {Monday
Tuesday}

d	w(R(d))	w($\neg R(d)$)
M	1	2
T	4	1

d	w(C(d))	w($\neg C(d)$)
M	3	5
T	6	2

Rain(M)	Cloudy(M)	Rain(T)	Cloudy(T)	Model?	Weight
T	T	T	T	Yes	$1 * 3 * 4 * 6 = 72$
T	F	T	T	No	0
F	T	T	T	Yes	$2 * 3 * 4 * 6 = 144$
F	F	T	T	Yes	$2 * 5 * 4 * 6 = 240$
T	T	T	F	No	0
T	F	T	F	No	0
F	T	T	F	No	0
F	F	T	F	No	0
T	T	F	T	Yes	$1 * 3 * 1 * 6 = 18$
T	F	F	T	No	0
F	T	F	T	Yes	$2 * 3 * 1 * 6 = 36$
F	F	F	T	Yes	$2 * 5 * 1 * 6 = 60$
T	T	F	F	Yes	$1 * 3 * 1 * 2 = 6$
T	F	F	F	No	0
F	T	F	F	Yes	$2 * 3 * 1 * 2 = 12$
F	F	F	F	Yes	$2 * 5 * 1 * 2 = 20$

$$+ \quad \quad + \quad \quad \quad \#SAT = 9 \quad WFOMC = 608$$

Gogate, V., & Domingos, P., Probabilistic Theorem Proving. Proc. UAI, 2012.

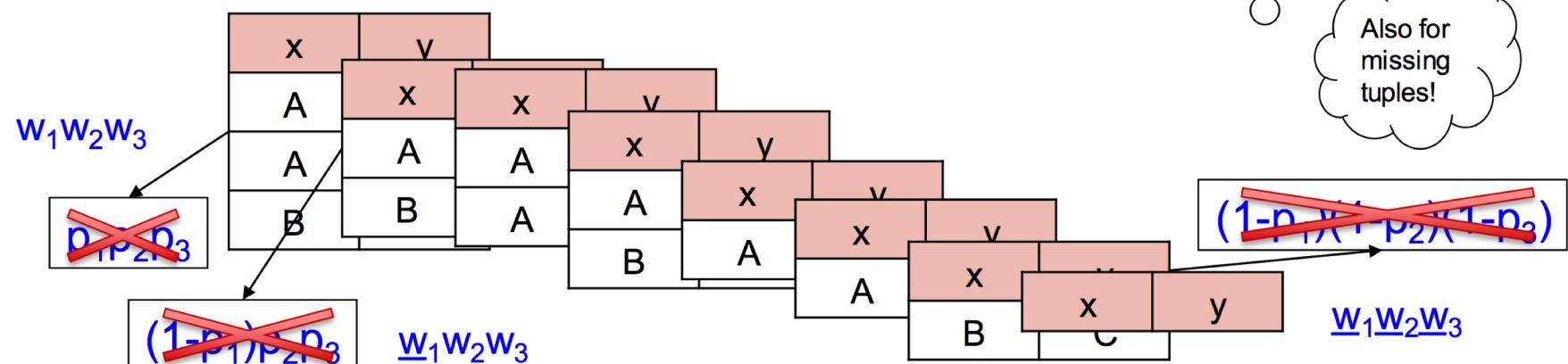
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Lifted probabilistic inference by first-order knowledge compilation.
In Proc.IJCAI-11, pp. 2178-2185, 2011.

From Probabilities to Weights

Friend		
x	y	P
A	B	p_1
A	C	p_2
B	C	p_3



x	y	w(Friend(x,y))	w(\neg Friend(x,y))
A	B	$w_1 = p_1$	$w_1 = 1-p_1$
A	C	$w_2 = p_2$	$w_2 = 1-p_2$
B	C	$w_3 = p_3$	$w_3 = 1-p_3$
A	A	$w_4 = 0$	$w_4 = 1$
A	C	$w_5 = 0$	$w_5 = 1$
	



Discussion

- Simple idea: replace $p, 1-p$ with w, \underline{w}
 - Weights, not necessarily probabilities
- Query answering by WFOMC
 - For obtaining probabilities:
Divide world weight by $Z = \text{sum of all world weights}$

Query Answering Problem

Given

MLN:

- 0.7 $\text{Actor}(a) \Rightarrow \neg \text{Director}(a)$
- 1.2 $\text{Director}(a) \Rightarrow \neg \text{WorkedFor}(a,b)$
- 1.4 $\text{InMovie}(m,a) \wedge \text{WorkedFor}(a,b) \Rightarrow \text{InMovie}(m,b)$

Database tuples (if missing w=1)

Actor:

Name	w
Brando	2.9
Cruise	3.8
Coppola	1.1

WorkedFor:

Actor	Director	w
Brando	Coppola	2.5
Coppola	Brando	0.2
Cruise	Coppola	1.7

Compute

$$P(\text{InMovie}(\text{GodFather}, \text{Brando})) = ??$$

Z Computation

1. Formula Δ

If all MLN constraints are hard: $\Delta = \bigwedge_{(\infty, \Gamma(x)) \in \text{MLN}} (\forall x \ \Gamma(x))$

If $(w_i, \Gamma_i(x))$ is a soft MLN constraint, then:

- a) Remove $(w_i, \Gamma_i(x))$ from the MLN
- b) Add new probabilistic relation $F_i(x)$
- c) Add hard constraint $(\infty, \forall x (F_i(x) \leftrightarrow \Gamma_i(x)))$

2. Weight function $w(.)$

For all constants A , relations F_i ,

set $w(F_i(A)) = \exp(w_i)$, $w(\neg F_i(A)) = 1$

Theorem: $Z = \text{WFOMC}(\Delta)$

Van den Broeck, G., Meert, W., & Darwiche, A., Skolemization for weighted first-order model counting. In Proc. KR-13, 2013.

Jha, A., & Suciu, D., Probabilistic databases with MarkoViews. Proceedings of the VLDB Endowment, 5(11), 1160-1171, 2012.

Example

1. Formula Δ

$\infty \text{ Smoker}(x) \Rightarrow \text{Person}(x)$

3.75 $\text{Smoker}(x) \wedge \text{Friend}(x,y) \Rightarrow \text{Smoker}(y)$

$\Delta = \forall x (\text{Smoker}(x) \Rightarrow \text{Person}(x))$
 $\wedge \forall x \forall y (F(x,y) \Leftrightarrow [\text{Smoker}(x) \wedge \text{Friend}(x,y) \Rightarrow \text{Smoker}(y)])$

2. Weight function $w(.)$

F			
x	y	$w(F(x,y))$	$w(\neg F(x,y))$
A	A	$\exp(3.75)$	1
A	B	$\exp(3.75)$	1
A	C	$\exp(3.75)$	1
B	A	$\exp(3.75)$	1
	

Note: if no tables given
for Smoker, Person, etc,
(i.e. no evidence)
then set their $w = \underline{w} = 1$

$Z = \text{WFOMC}(\Delta)$

Knowledge Compilation for Counting

- Main idea: convert Δ into a different “form” from which one can easily read off the solution count (and many other quantities of interest)
[Darwiche & Marquis 2002]
- d-DNNF: Deterministic, Decomposable Negation Normal Form
 - Think of the formula as a directed acyclic graph (DAG)
 - Negations allowed only at the leaves (NNF)
 - Children of AND node don’t share any variables (different “components”)
 - Children of OR node don’t share any solutions
 - can add the counts
 - can multiply the counts
- Once converted to d-DNNF, **can answer many queries in linear time**
 - Satisfiability, tautology, logical equivalence, solution counts, ...
 - Any query that a BDD could answer

Compilation to d-DNNF

- “Domain-liftable” FO formula

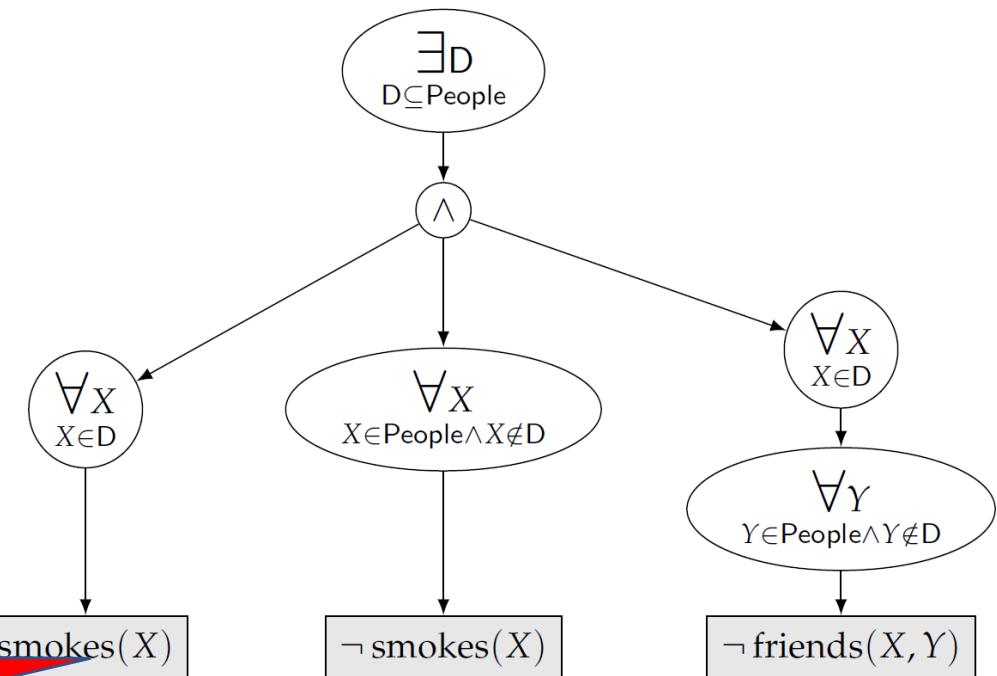
$\forall X, Y \in \text{People}$

$\text{smokes}(X) \wedge \text{friends}(X, Y)$

$\Rightarrow \text{smokes}(Y)$

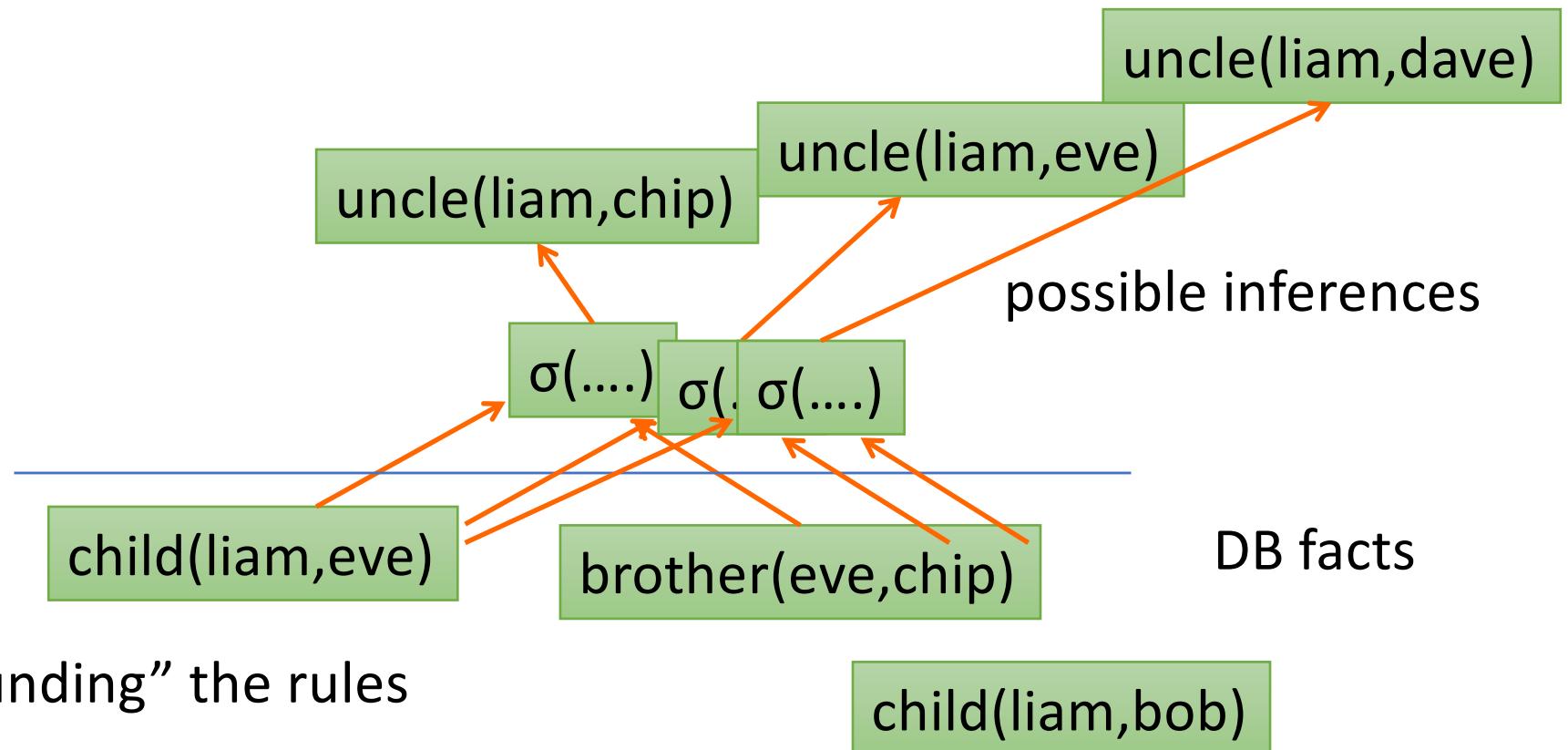
- Probability of a query depends only on the **size(s) of the domain** (s), a **weight function** for the first-order predicates, and the **weighted model count** over the FO d-DNNF.

d-DNNF form of Δ
can grow large



TensorLog [Cohen & Yang 17]

```
. uncle(X,Y) :- child(X,W), brother(W,Y) .  
. uncle(X,Y) :- aunt(X,W), husband(W,Y) .  
. status(X,tired) :- child(W,X), infant(W) .
```



“Grounding” the rules

William W. Cohen and Fan Yang, TensorLog: Deep Learning Meets Probabilistic Databases in arxiv.org 1707.05390, 2017

Explicit grounding not scalable

```
:uncle(X,Y) :- child(X,W), brother(W,Y).  
:uncle(X,Y) :- aunt(X,W), husband(W,Y).  
status(X,tired) :- child(W,X), infant(W).
```

Example: inferring family relations like “uncle”

- N people
- N^2 possible “uncle” inferences
- $N = 2 \text{ billion} \rightarrow N^2 = 4 \text{ quintillion}$
- $N = 1 \text{ million} \rightarrow N^2 = 1 \text{ trillion}$

A KB with 1M entities is *small*

Key Question: How to reason?

```
uncle(X,Y) :- child(X,W), brother(W,Y).  
uncle(X,Y) :- aunt(X,W), husband(W,Y).  
status(X,tired) :- child(W,X), infant(W).
```

Example: inferring family relations like “uncle”

- N people
 - N^2 possible “uncle” facts
 - ~~$N = 1 \text{ million} \rightarrow N^2 = 1 \text{ trillion}$~~

Functions f_i
to be learned
from data

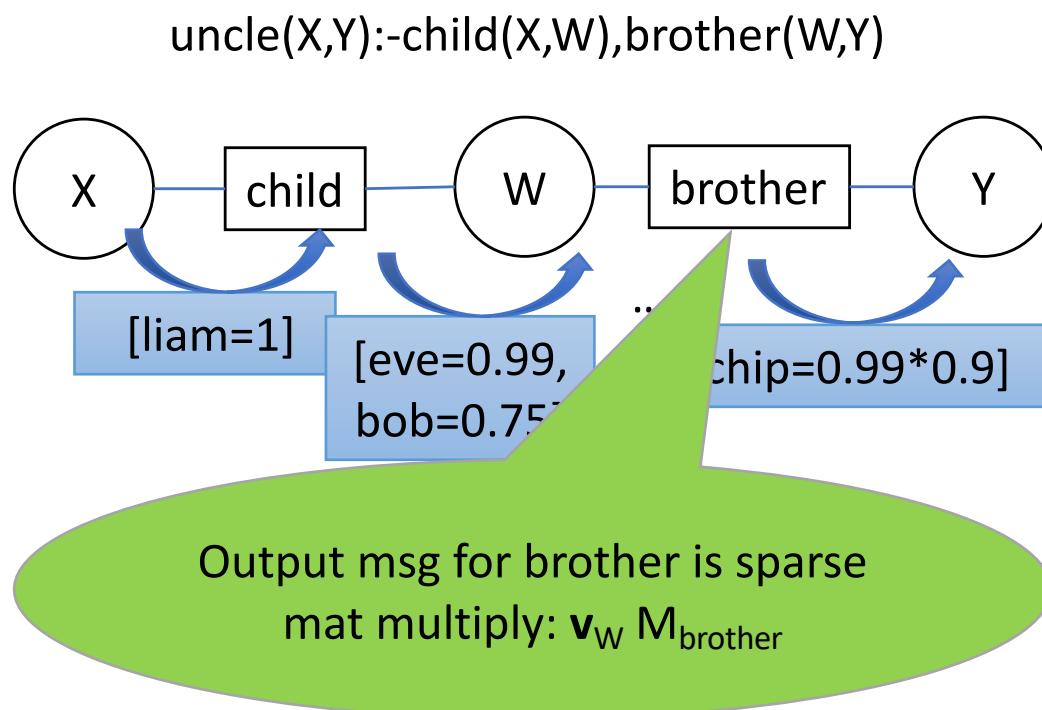
The diagram illustrates a mapping from input x to output Y . The input x is labeled "one-hot vectors" and is shown as a blue arrow pointing upwards from the bottom. The output Y is labeled "f₁(x)" and is shown as a blue arrow pointing to the right. The background is light blue.

x is the uncle
 $f_2(x) = Y$

vectors encoding
weighted set of DB instances

TensorLog: Approximation by Belief Propagation

Query: uncle(liam, Y) ?



child(liam,eve),0.99	infant(liam),0.7
child(dave,eve),0.99	infant(dave),0.1
child(liam,bob),0.75	aunt(joe,eve),0.9
husband(eve,bob),0.9	brother(eve,chip),0.9

General case for $p(c,Y)$:

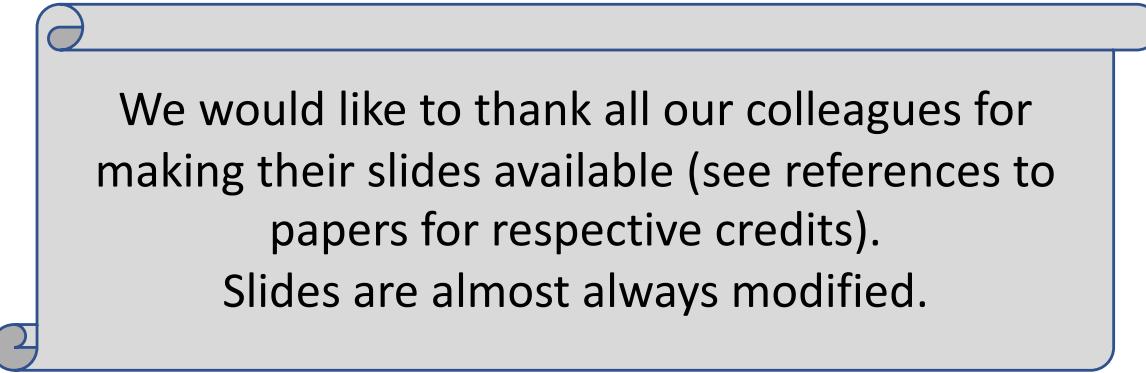
- initialize the evidence variable X to a one-hot vector for c
- wait for BP to converge
- read off the message \mathbf{y} that would be sent from the output variable Y .
 - un-normalized prob
 - $\mathbf{y}[d]$ is the **weighted number of proofs supporting $p(c,d)$ using this clause**

Wrap-up Statistical Relational AI

- Probabilistic relational logics
 - Overview
 - Semantics
 - Inference problems
- Dealing with scalability issues (avoiding grounding)
 - Reduce expressivity (liftable queries)
 - Knowledge compilation (WFOMC)
 - Approximation (BP)

Next: Lifted Exact Inference

Acknowledgments



We would like to thank all our colleagues for making their slides available (see references to papers for respective credits). Slides are almost always modified.